

# Type Inference in Object-Oriented Languages with Classes for Linguistic Engineering

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## Context and Goal

- Context  
⇒ Linguistic Engineering
- Goal  
⇒ Type Inference in Object-Oriented Languages with Classes

## A General Example from Botany (1/2)

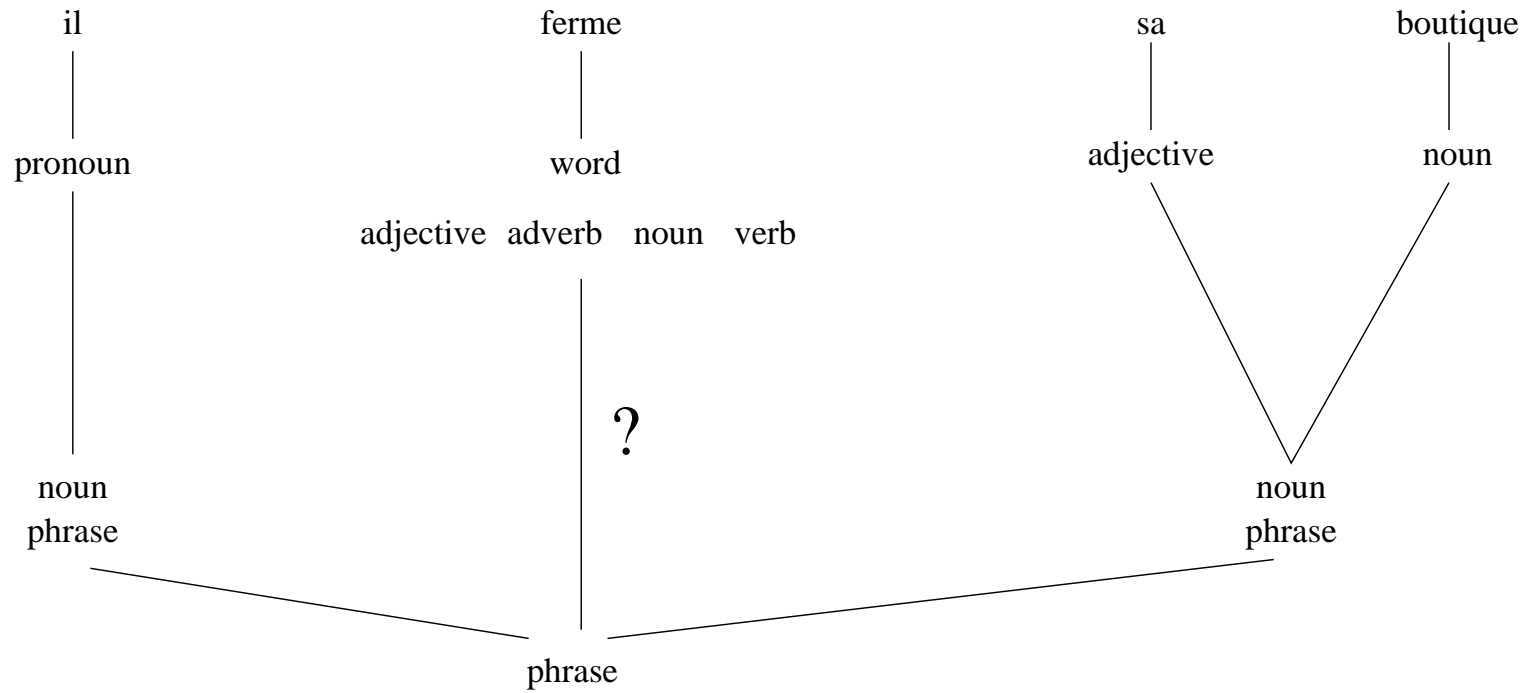
- We can see a plant (*object*).  
     $\implies$  its species (*type*) are Plant
- It has green leaves (*value* and *attribute*)  
     $\implies$  its species are Chlorophyllian (*subtype* of Plant)
- and flowers (*attribute*).  
     $\implies$  its species are Phanerogam (*subtype* of Chlorophyllian)

## A General Example from Botany (2/2)

At the end of the analysis, we can deduce:

- either the exact species of the plant,
- or a set of possible species for the plant,
- or a species discovery or an analysis error.

# An Example from Linguistics



*he closes his shop*

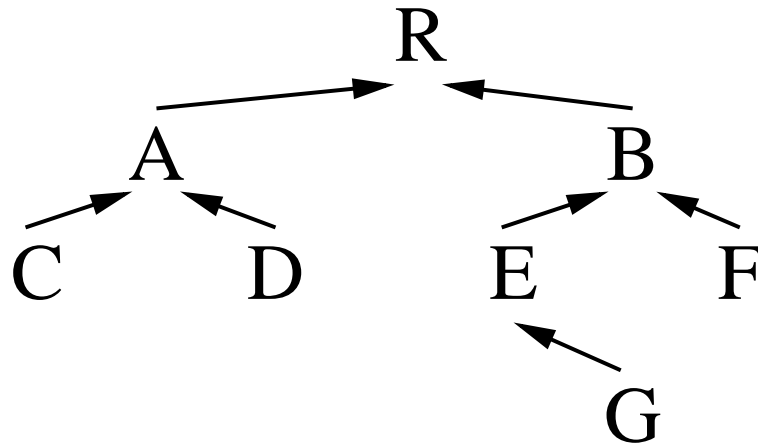
## About Object-Orientation

- **Reification**
- **Encapsulation**
- **Specialization**
  - Class and Type
  - Primitives: Attribute and Value, Method and Behavior
  - Object and Instance
  - Utilization Links: Aggregation and Composition
  - Importation Link : Inheritance
  - Compatibility
  - Migration

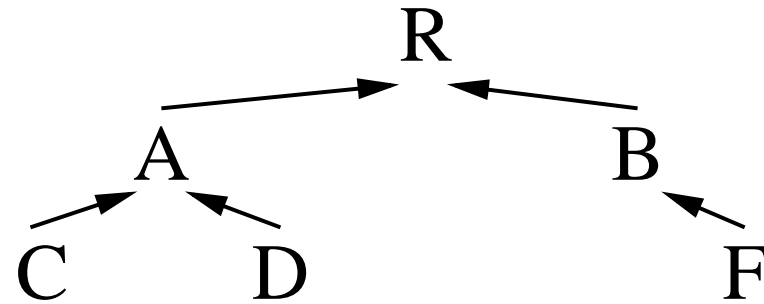
## First Steps to a Solution

1. Instanciation with the most general type
2. Data acquisition
3. Type refining (migration)
4. Results interpretation

# Our Type System: Typological Domain (1/2)



x is not of type E





## Our Type System: Typological Domain (2/2)

### Results Interpretation:

- The final typological domain
  - contains only one type, or
  - is empty, or
  - contains several types.

## Constraint-Controlled Migration

$o.p = 3$

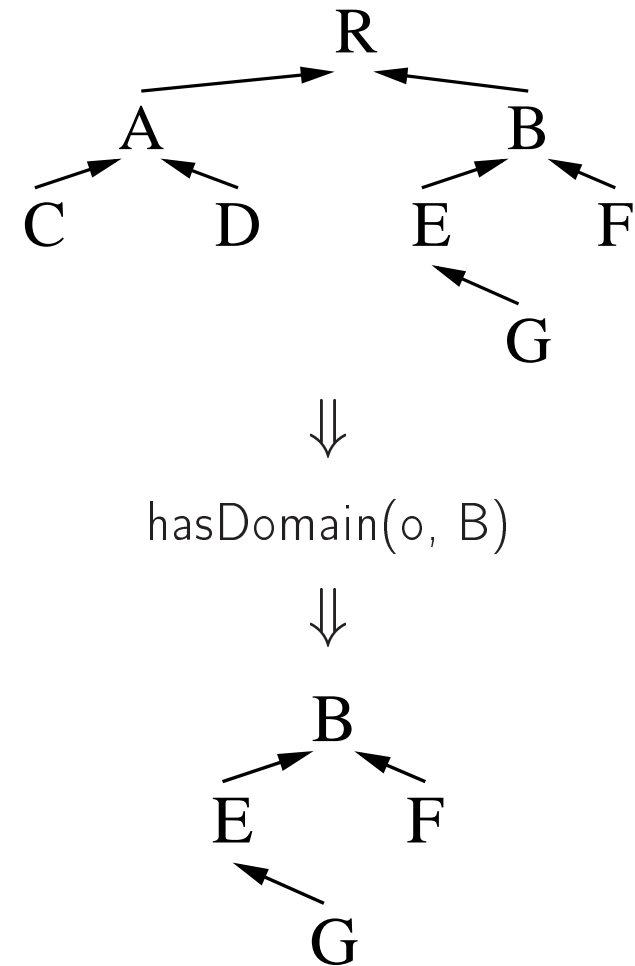
$\Rightarrow o$  has a primitive named  $p$

$\Rightarrow p$  of  $o$  is an Integer

## hasDomain Constraint

$\text{hasDomain}(o, D)$

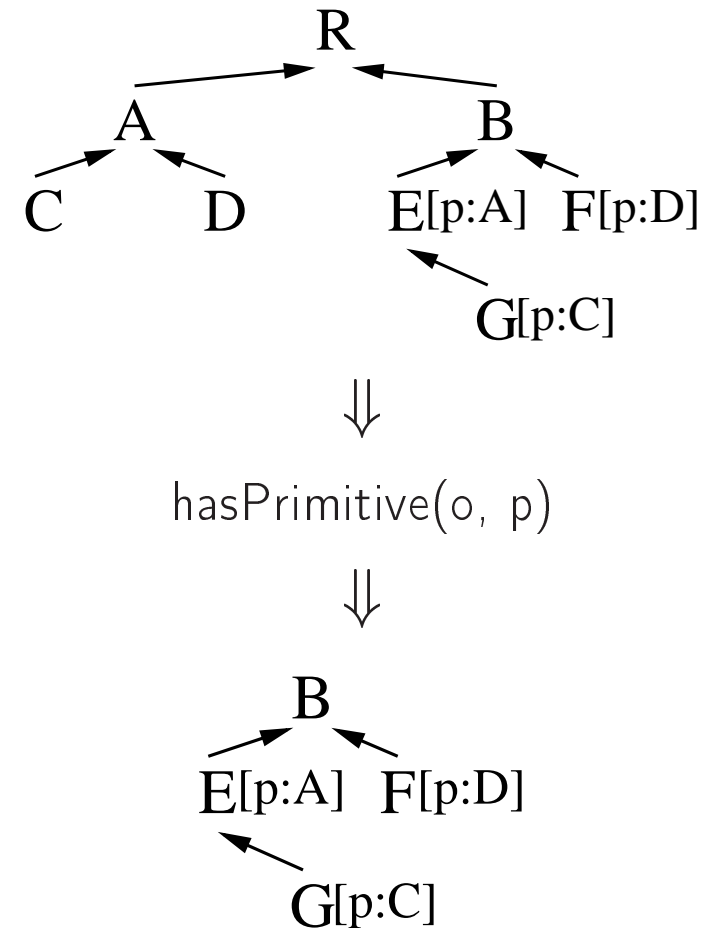
1. The new typological domain (TD) of  $o$  is the intersection of its current TD with  $D$ .
2. This constraint is recursively propagated over all the objects which reference and are referenced by the new  $o$ .



## hasPrimitive Constraint

$\text{hasPrimitive}(o, p)$

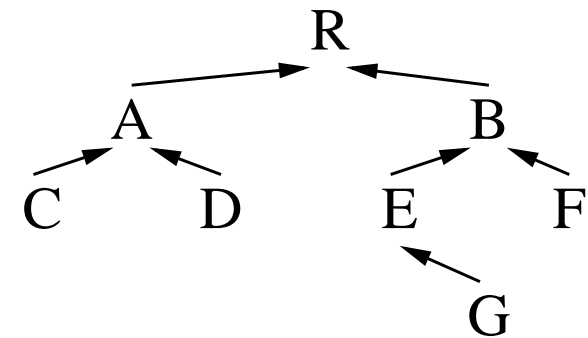
1. If  $o$  hasn't a  $p$  primitive and there is at least one  $p$  primitive in the current TD of  $o$ , then a  $p$  primitive is added to  $o$ . The new TD of  $o$  and  $p$  are refined.
2. This constraint is recursively propagated over all the objects which reference and are referenced by the new  $o$  and  $p$ .



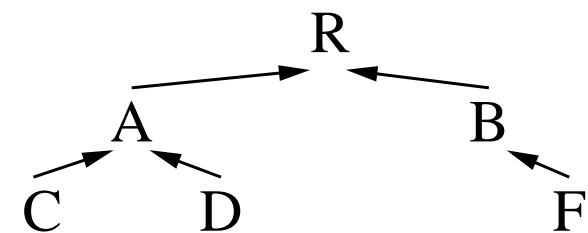
## hasNotType Constraint

$\text{hasNotType}(o, T)$

1. The new TD of  $o$  is the current TD of  $o$  minus the TD descended from  $T$ .
2. This constraint is recursively propagated over all the objects which reference and are referenced by the new  $o$ .



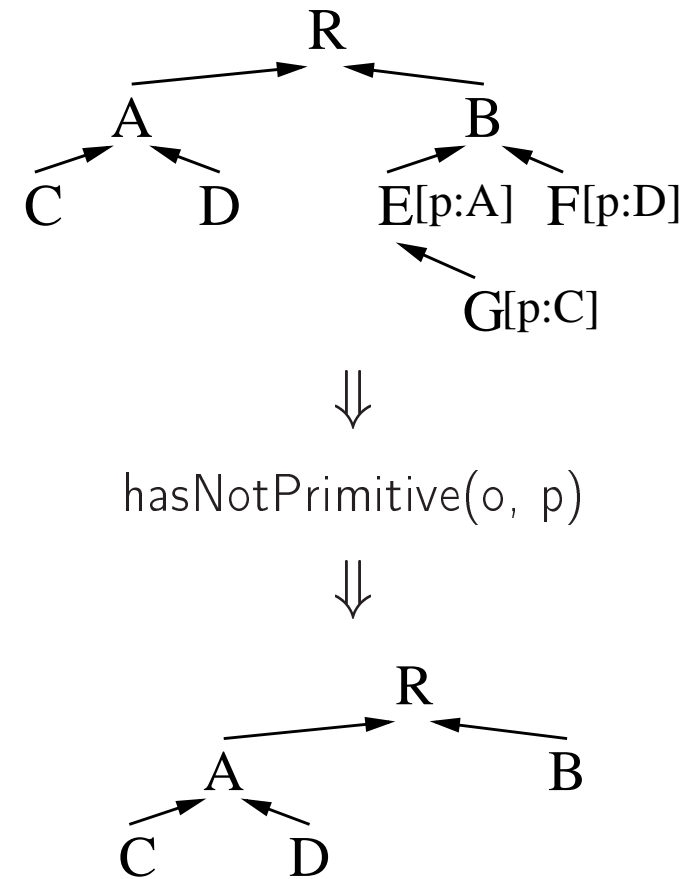
$\text{hasNotType}(o, E)$



## hasNotPrimitive Constraint

$\text{hasNotPrimitive}(o, p)$

1. The new TD of  $o$  is the smallest subtree (of the current TD of  $o$ ) which doesn't contain  $p$ .
2. This constraint is recursively propagated over all the objects which reference and are referenced by the new  $o$ .



## About Language Operators

creation of $o$ with type $T$	$allocate(o)$ $hasDomain(o, domainOf(T))$ $initialize(o)$
access to $p$ primitive of $o$	$hasPrimitive(o, p)$ $o.p$
assignment of $o'$ to $o$	$hasDomain(o, domainOf(o'))$ $o = o'$

## Conclusion and Prospects

Our current work:

- extends the type system of Object-Oriented Languages with Classes,
- provides a type inference mechanism,
- could be adapted to handle exceptions, and
- will be prototyped in CLOS then implemented as an extension of Java.