

M2 Complex Systems – Complex Networks

Lecture 13

Complex networks as almost structured graphs

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Modelling static networks

MODEL = RANDOM GENERATION OF SYNTHETIC NETWORKS

For simulating:

- phenomena
- algorithms
- protocols

In order to:

- design
- test
- predict
- better understand

Q: Do Internet protocols still work if Internet is 10 times larger ?



Generate a synthetic network and simulate

Modelling static networks

MODEL = RANDOM GENERATION OF SYNTHETIC NETWORKS

4 classic properties:

- Low global density
- Short distances
- Heterogeneous degrees
- High local density

Modelling static networks

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Modelling static networks

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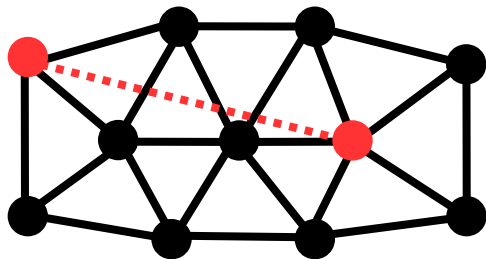
Modelling static networks

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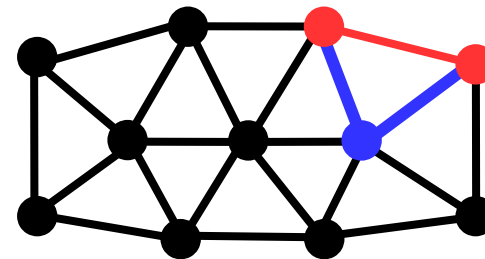
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- High local density → **problem**

proba ???



global density

proba ???



local density

Modelling static networks

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Big challenge: Generate networks having these 4 properties



low global density



short distances



heterogeneous degrees



high local density

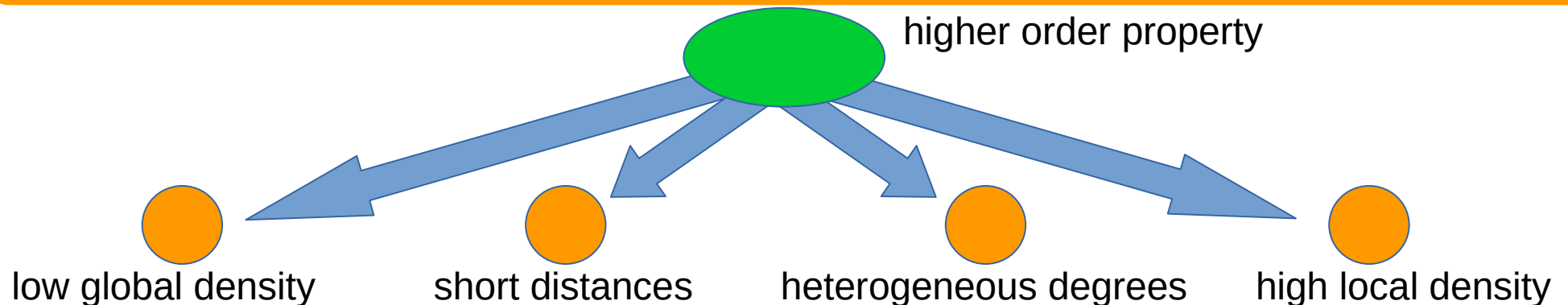
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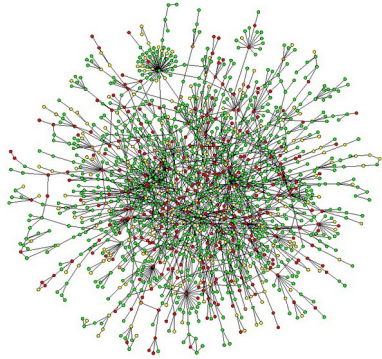
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Big challenge: Generate networks having these 4 properties



Idea: obtain these properties as a consequence of a higher order property

Almost structured graphs



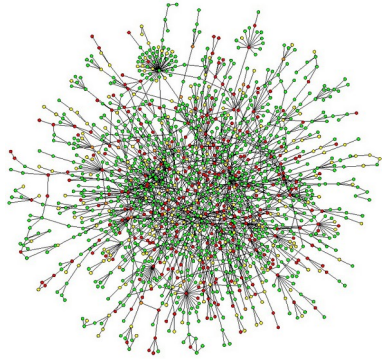
- loosely constrained

→ randomness

- strongly impacted by their context

→ structure

Almost structured graphs



■ loosely constrained

→ randomness

■ strongly impacted by their context

→ structure

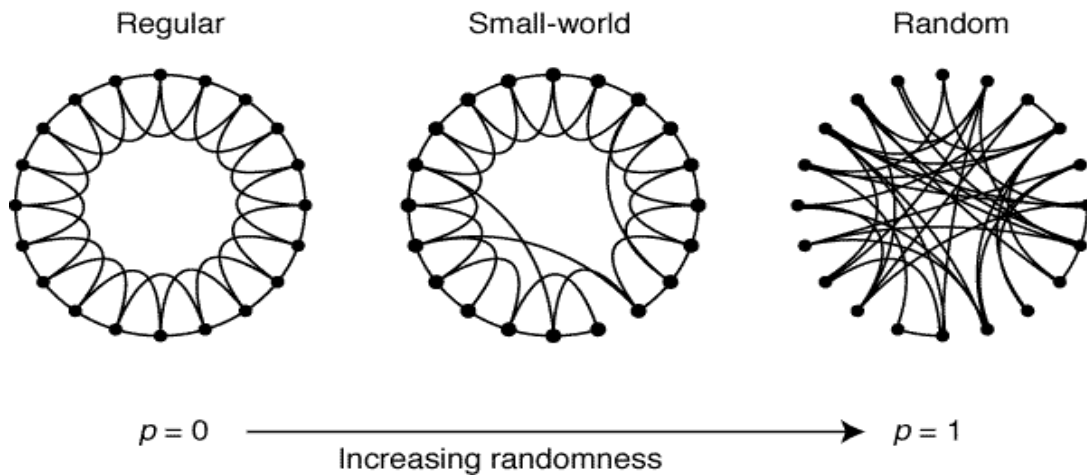
Complex networks = **structure** + **randomness**

[Watts & Strogatz 1998]

High local density

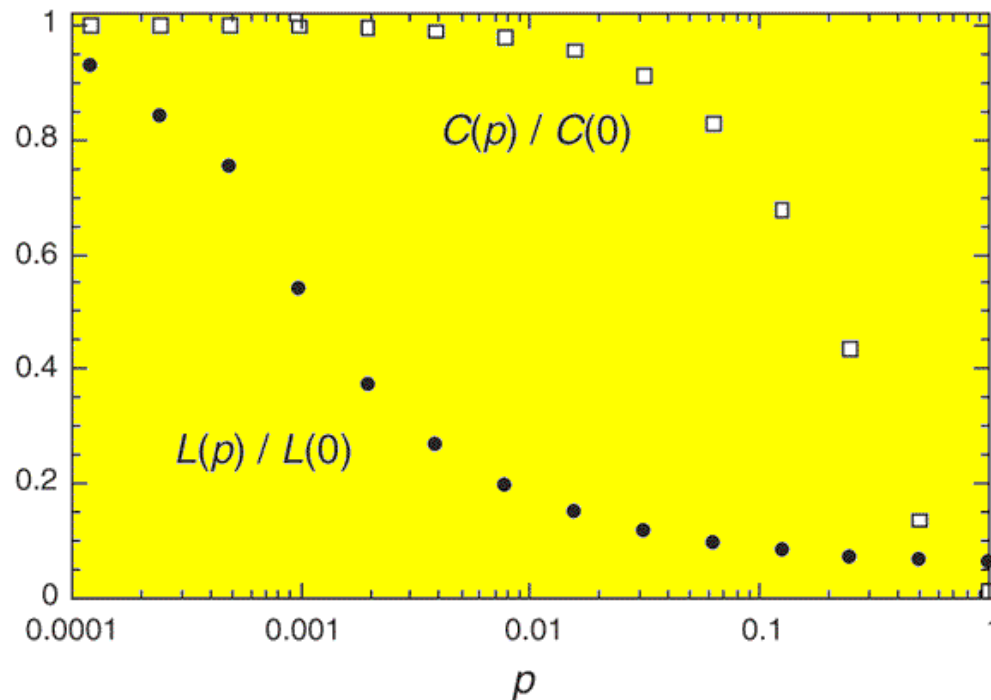
Short distances

Watts & Strogatz model



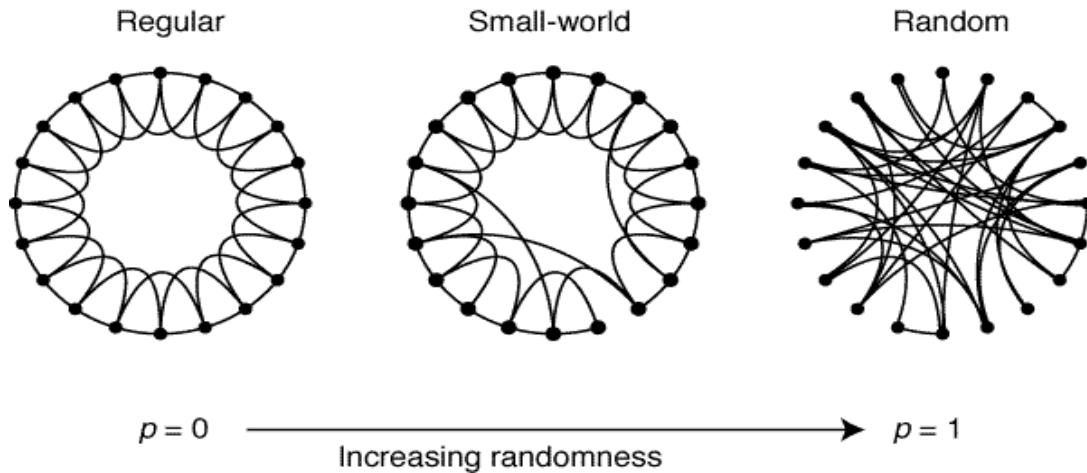
Regular lattice with n nodes
 k^{th} power of the cycle, $k \ll n$

Second endpoint of each edge
is rewired with probability p



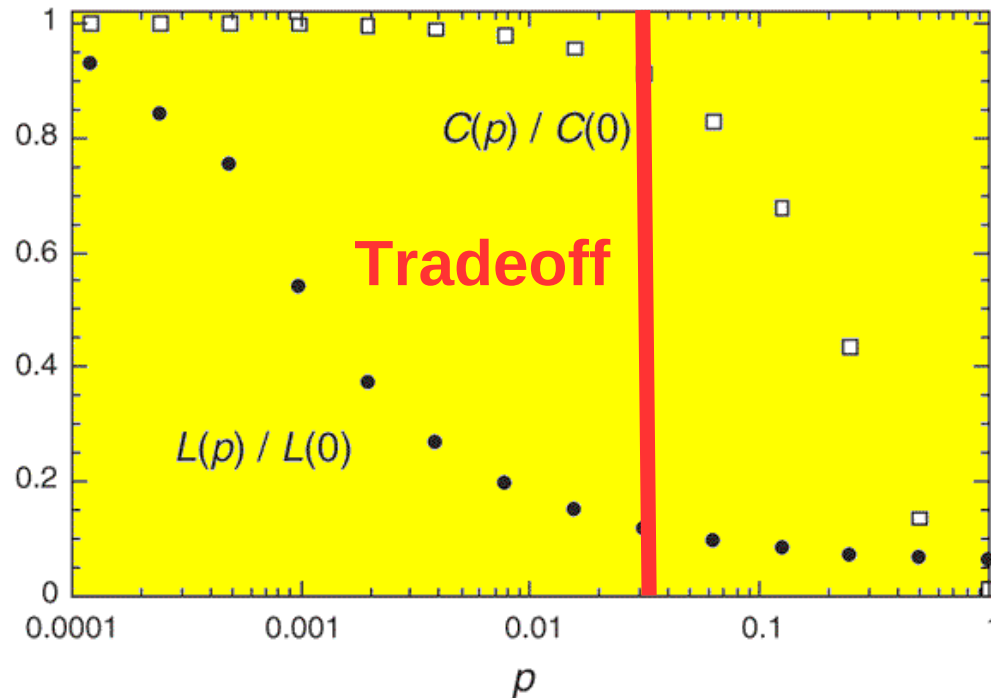
Clustering $C(p)$
VS
average distance $L(p)$
as p increases

Watts & Strogatz model



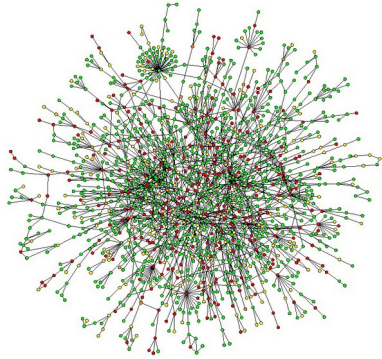
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Almost structured graphs



loosely constrained

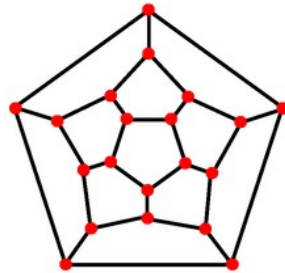
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strongly impacted by their context

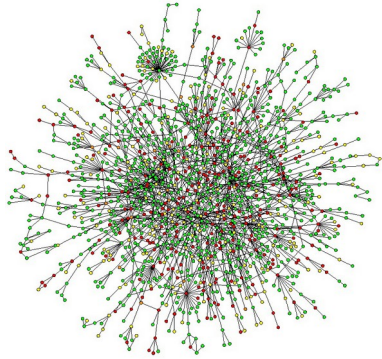
→ structure

Complex networks = structure + randomness

1 strongly structured



Almost structured graphs



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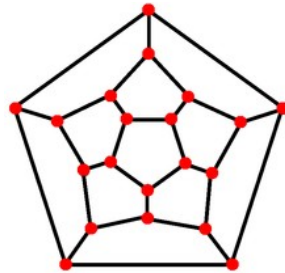
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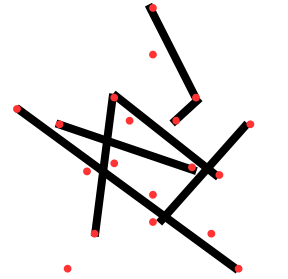
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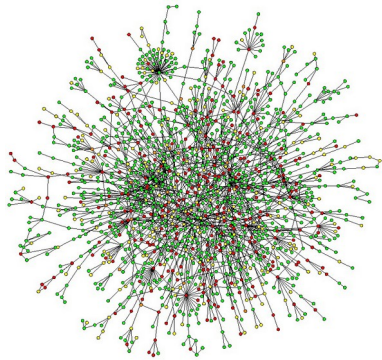
1 strongly structured



2 random modifications



Almost structured graphs



loosely constrained

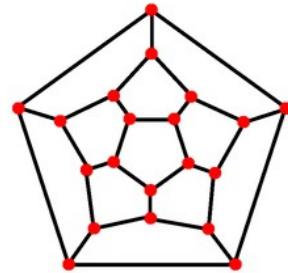
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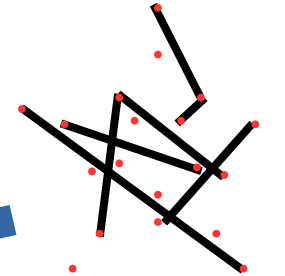
→ structure

Complex networks = structure + randomness

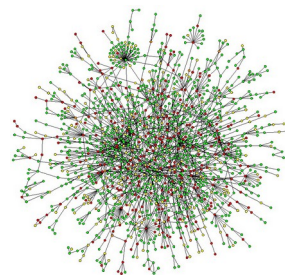
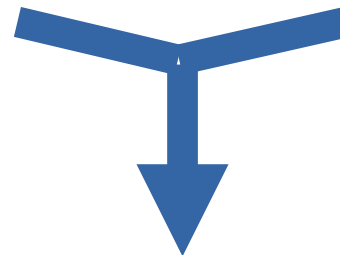
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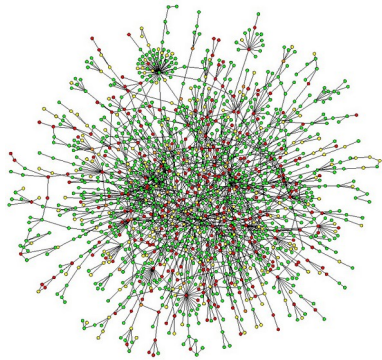
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+



Almost structured graphs



loosely constrained

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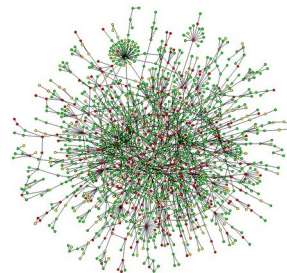
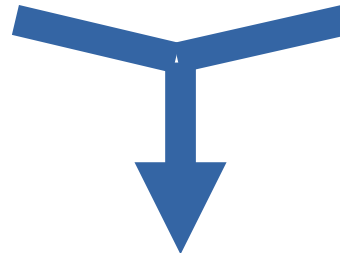
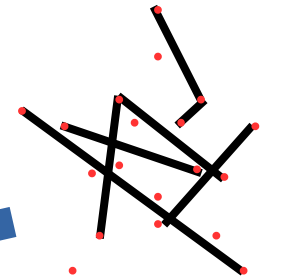
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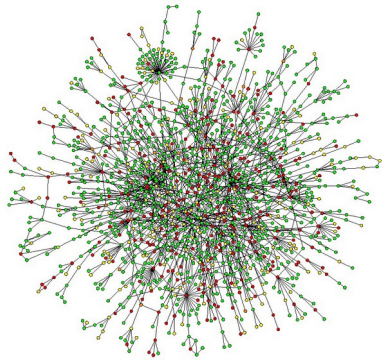
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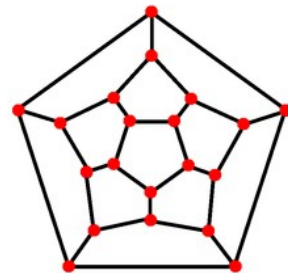
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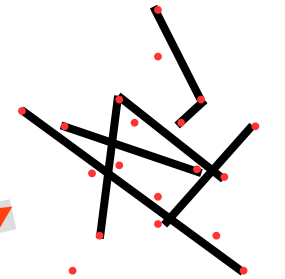
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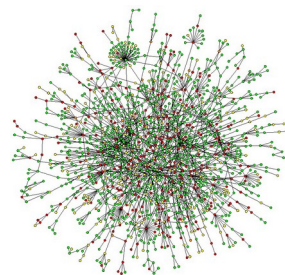
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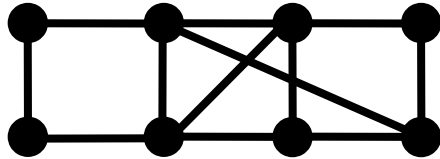
structure

noise



Graph editing algorithms

INPUT

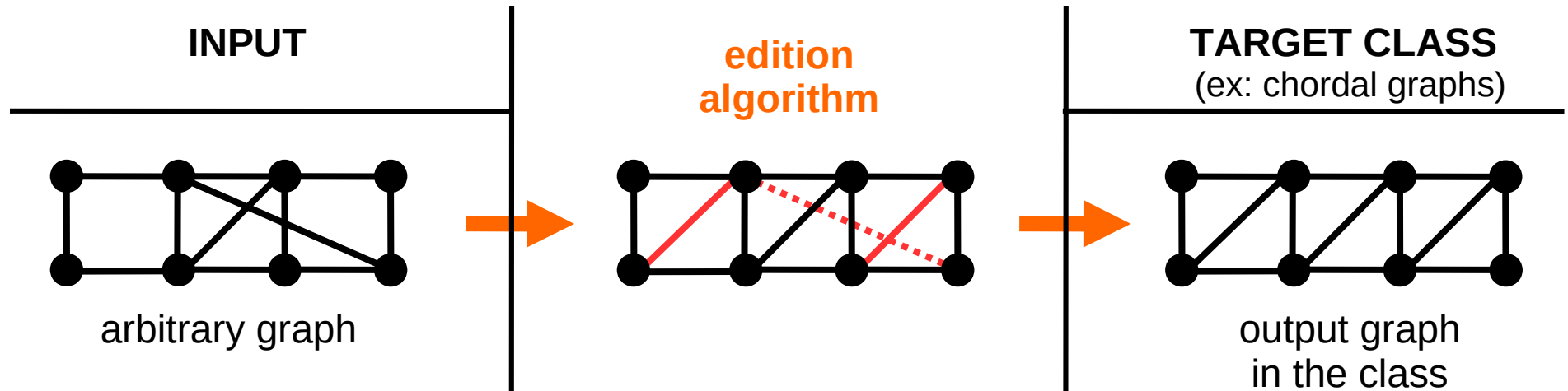


arbitrary graph

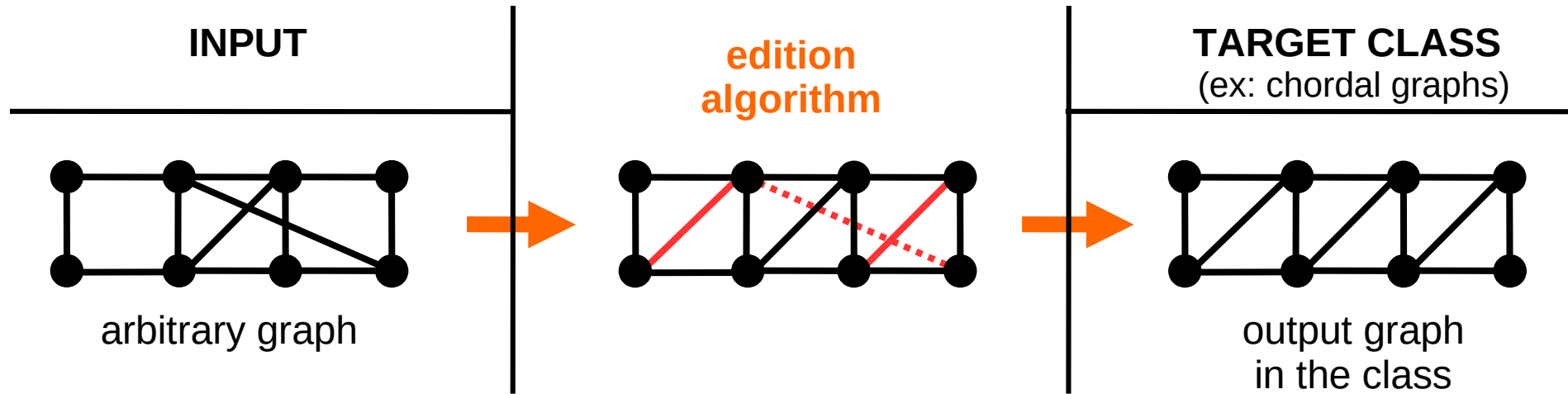
TARGET CLASS

(ex: chordal graphs)

Graph editing algorithms



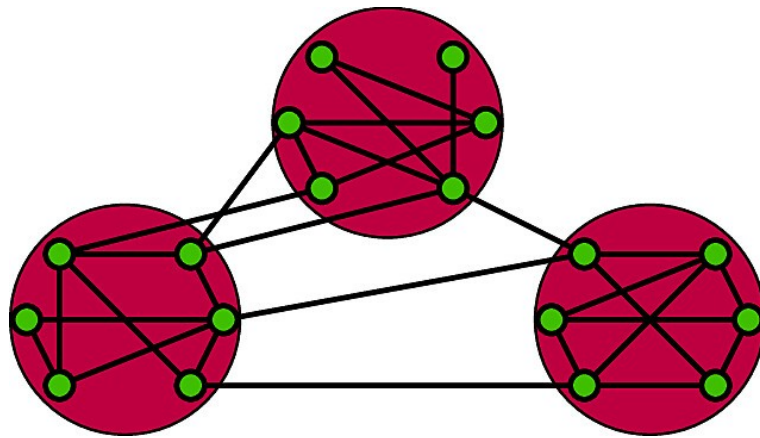
Graph editing algorithms



GOAL: perform as few modifications as possible

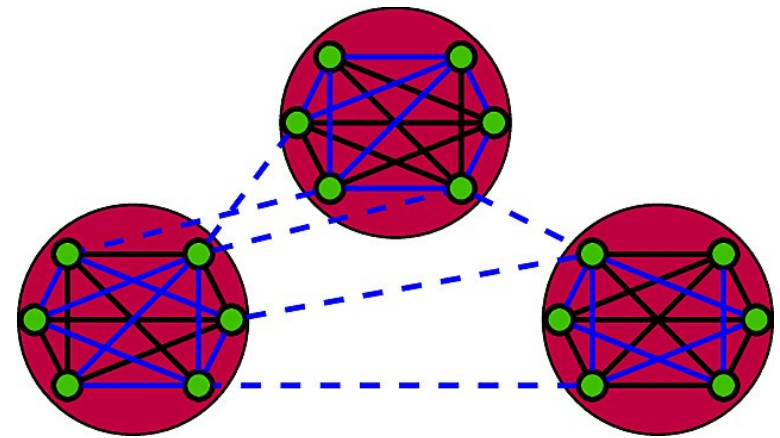
Graph editing algorithms

■ Community detection



Original network

EDITING

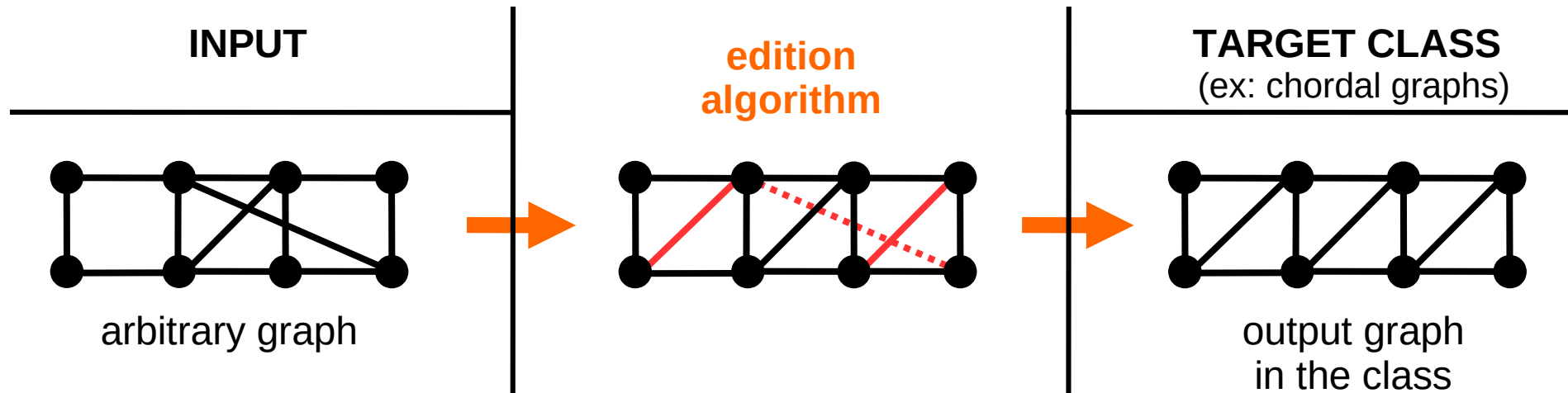


Resulting cluster graph

■ Degree anonymization

- Edit the graph so that all vertices have same degree

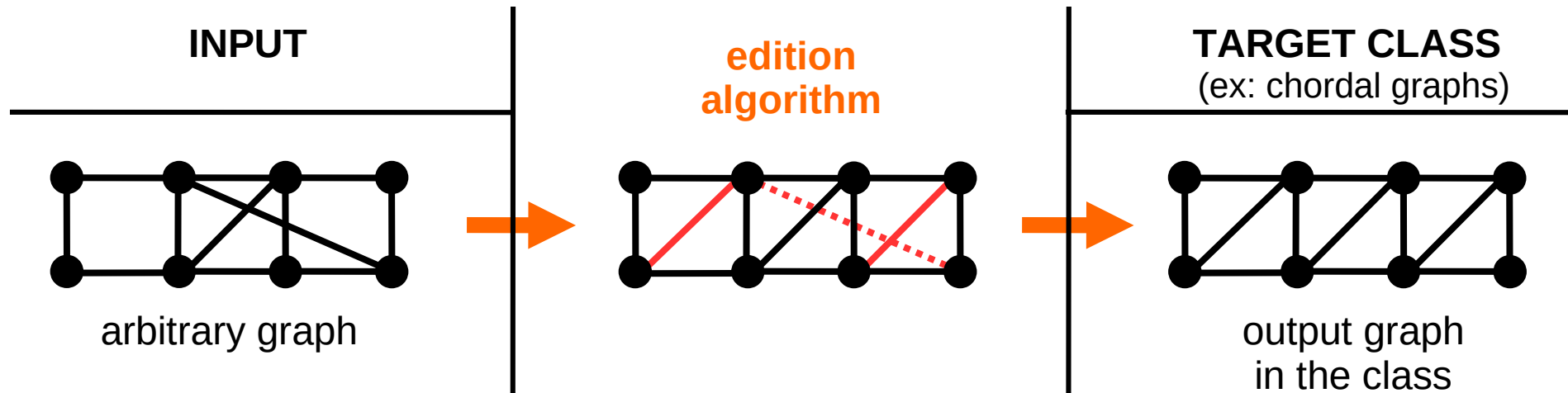
Graph editing algorithms



GOAL: perform as few modifications as possible

- Unfortunately: *minimum number* is **NP-hard** for most properties
 - Even when only one type of modifications is allowed (eg. only additions)
- Different approaches:
 - Restricted inputs
 - Exact exponential algorithms
 - Parameterized algorithms
 - Approximation algorithms
 - *Inclusion minimal modification*

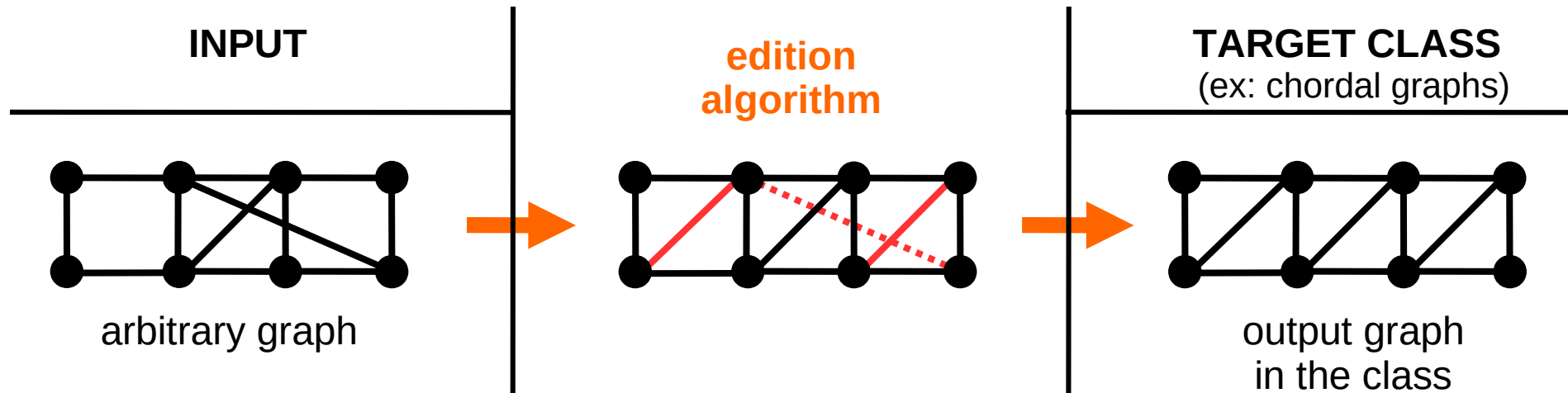
Graph editing algorithms



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Even when only one type of modifications is allowed (eg. only additions)
- Relaxation of the problem:
set of modifications *minimal for inclusion* → **polynomial time**

Graph editing algorithms



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Even when only one type of modifications is allowed (eg. only additions)

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each target class needs a **specific** algorithm !

Ex : interval graphs, permutation graphs, cographs

Results for some target classes

Completion:

- Interval completion : $O(n^2)$
1981, 2005, 2013
- Chordal completion : $O(nm)$
2006
- Trivially perfect completion : $O(n+m')$
2008
- Comparability completion : $O(n^2m)$
2008
- Split completion : $O(n+m')$
2009
- Cograph completion : $O(n+m')$
2010
- Permutation completion : $O(n^2)$
2015

Deletion:

- Planar deletion : $O(n+m)$
2006

Minimal cograph editing algorithms

Coedit : a tool for cograph editing

INPUT: an arbitrary graph

Computes either:

- a minimal cograph completion
- a minimal cograph deletion
- a minimal cograph editing

OUTPUT: the cotree of the cograph obtained

Input format:

# of vertices	n
degrees	$\begin{cases} u & d^\circ(u) \\ v & d^\circ(v) \\ \vdots & \end{cases}$
edges	$\begin{cases} u1 & v1 \\ u2 & v2 \\ \vdots & \end{cases}$

Output format:

# of nodes	n
Label of the root	l (=0 or 1)
# of children	$\begin{cases} u & \#child(u) \\ v & \#child(v) \\ \vdots & \end{cases}$
Edges of the tree	$\begin{cases} parent(u) & u \\ parent(v) & v \\ \vdots & \end{cases}$

- Written in C
- Sources available at <https://www.ii.uib.no/~christophec/coedit/>
- Under GNU GPL licence (can do whatever you want with it)

Algorithms

For completion

- An $O(n+m')$ algorithm with *minimum* at each incremental step
➔ improve heuristics
- An $O(n+m \log^2 n)$ algorithm
➔ almost linear in the size of the *input*

For editing

- An $O(n+m)$ algorithm with *minimum* at each incremental step

The vertex incremental approach : vertices are processed one by one



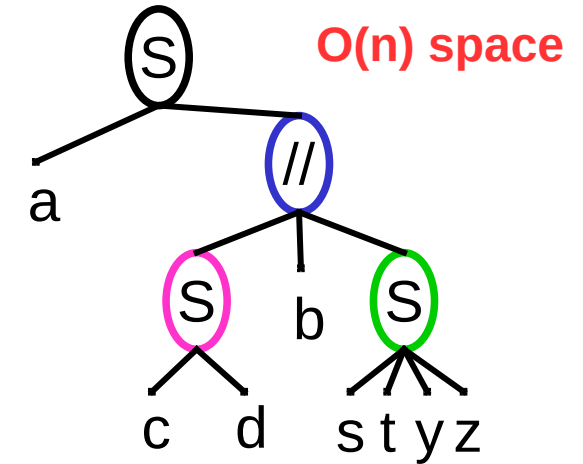
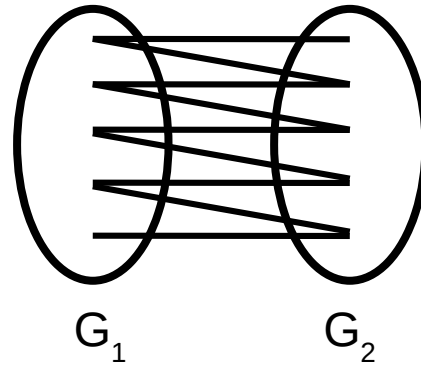
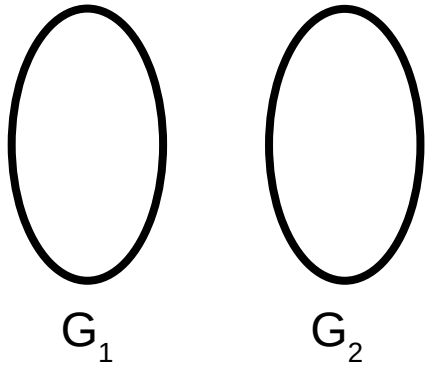
Cographs and incremental app.

Obtained from single vertices by using 2 operations:

cotree

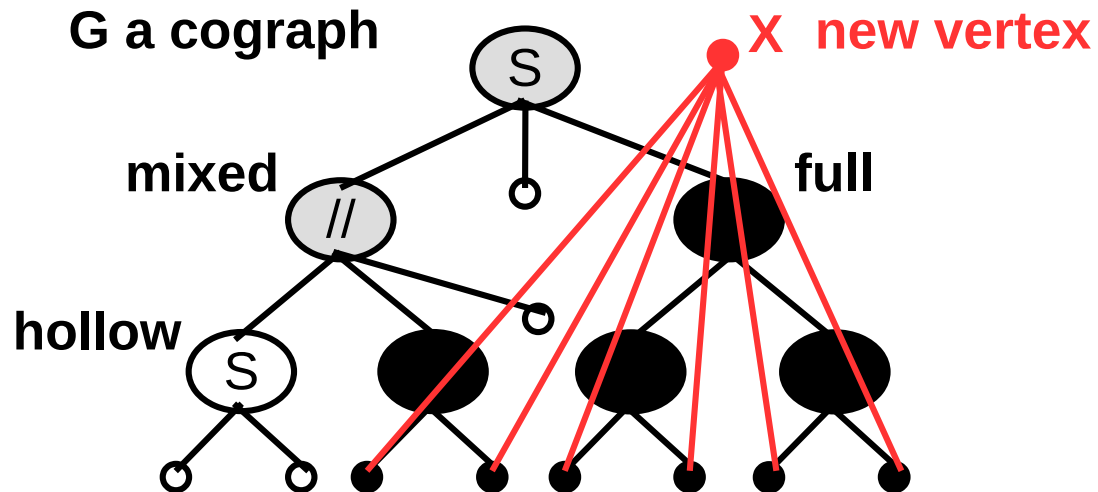
disjoint union
(//)

complete union
(S)



Incremental approach: a **cograph G** and **x** a new incoming vertex

$G+x$ is not a cograph and we want to add (and/or delete) edges incident to x so that $G+x$ become a cograph



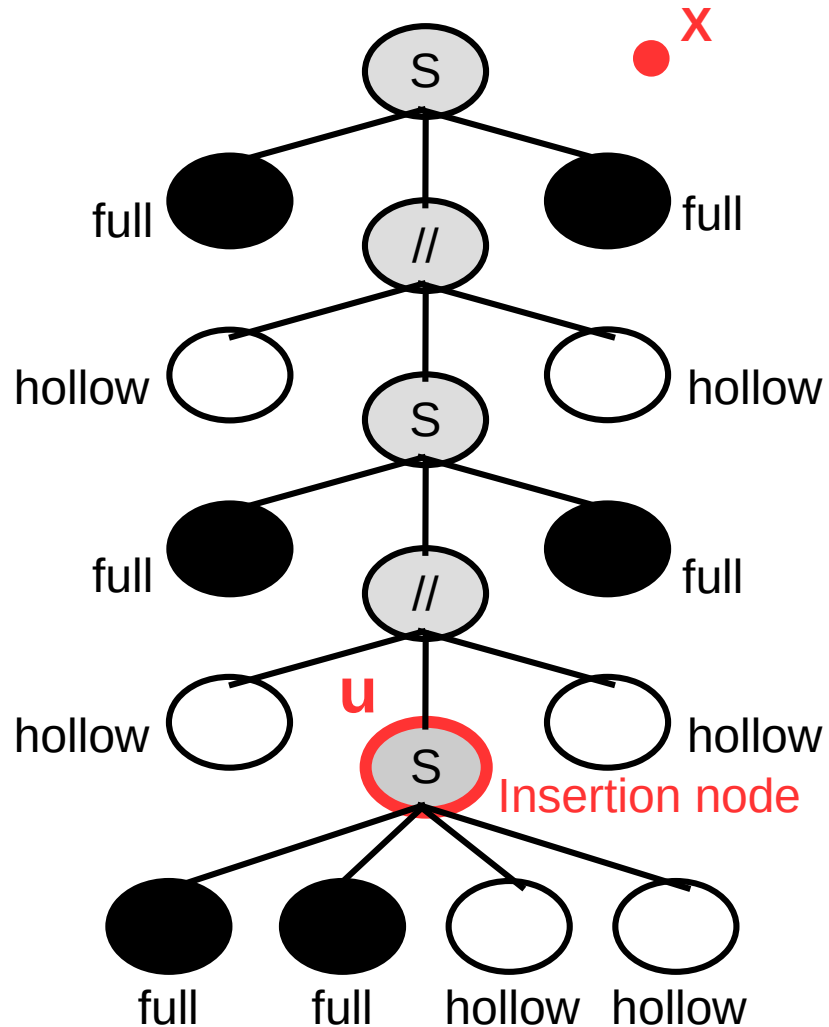
Completion algorithms

First algorithm: $O(n+m')$

A characterisation of cographs

[Corneil, Perl, Stewart 1981]

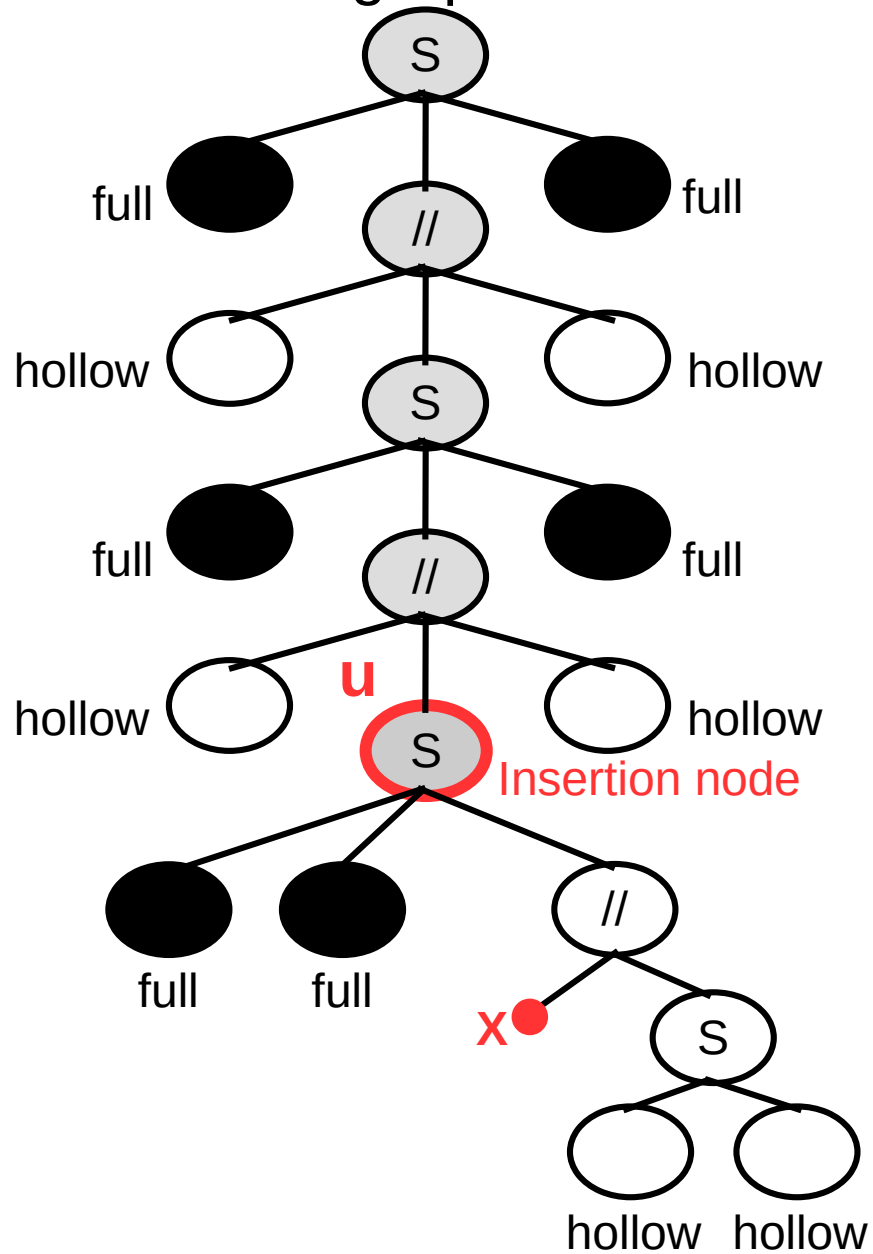
$G+x$ is a cograph iff there exists a node u st.:



A characterisation of cographs

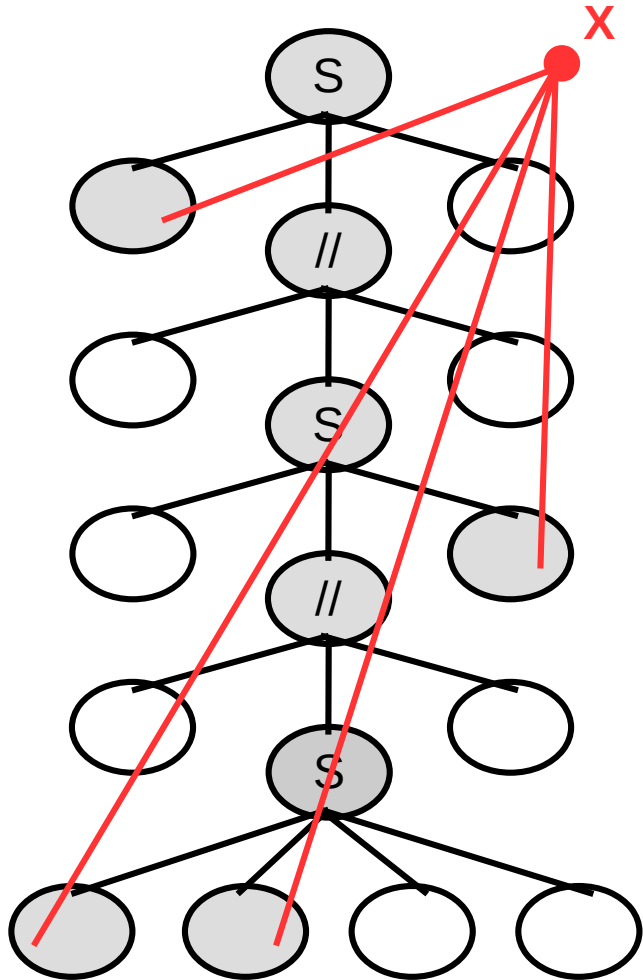
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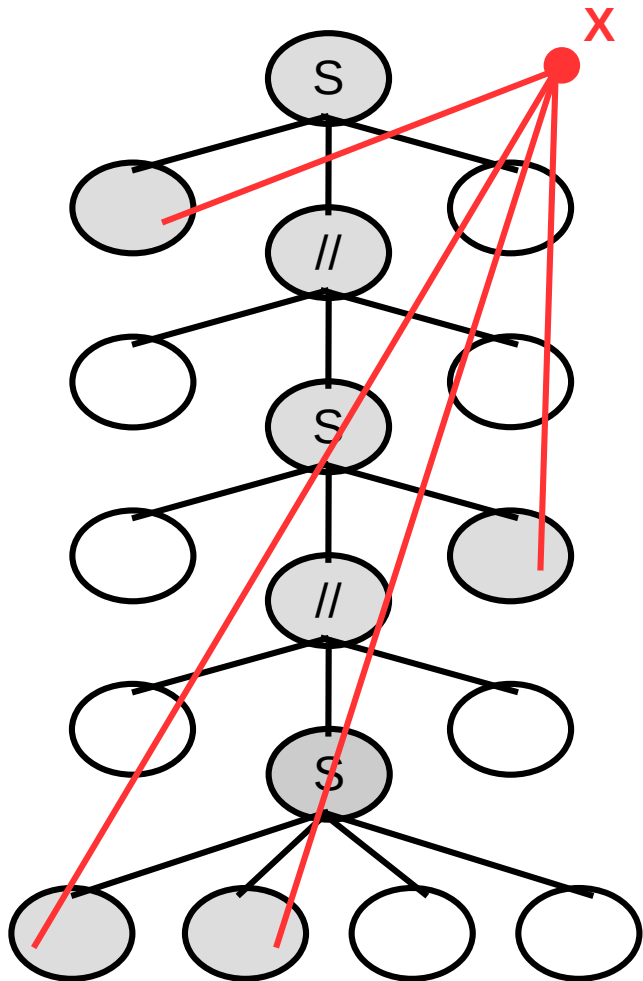
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In our algorithm : $G+x$ is not a cograph



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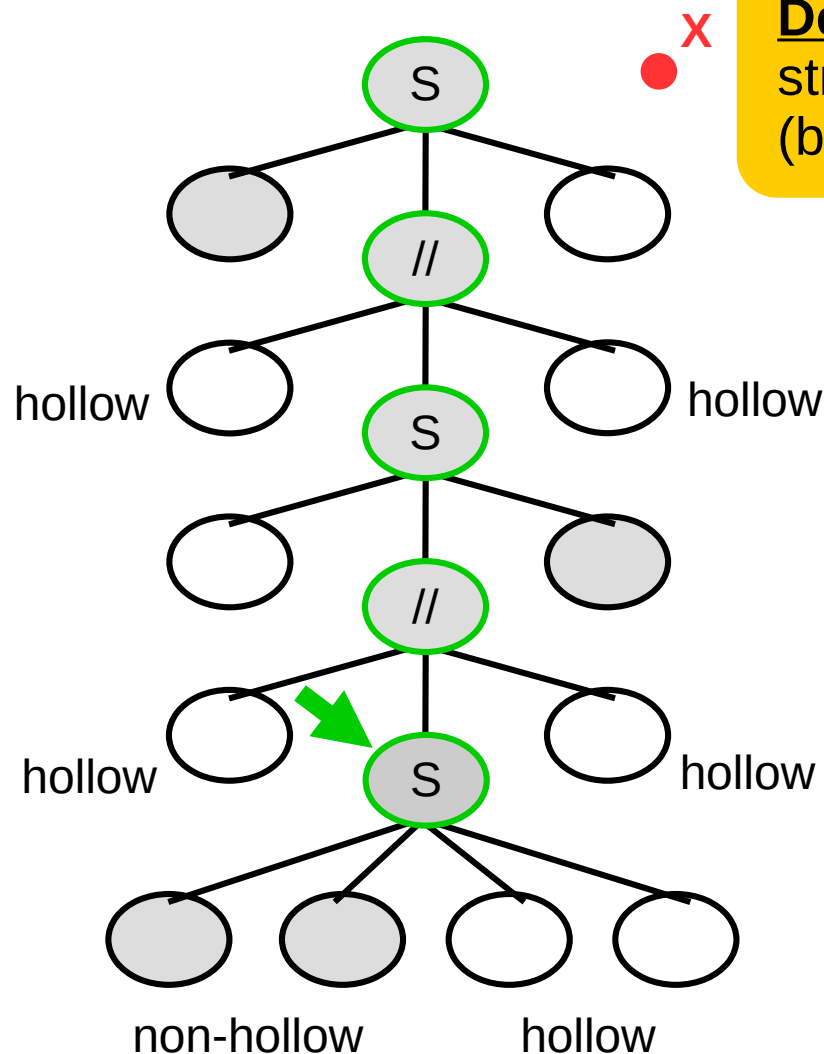
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Choose one node u for which you make the situation of the [CPS 81]'s theorem happen

Eligible nodes

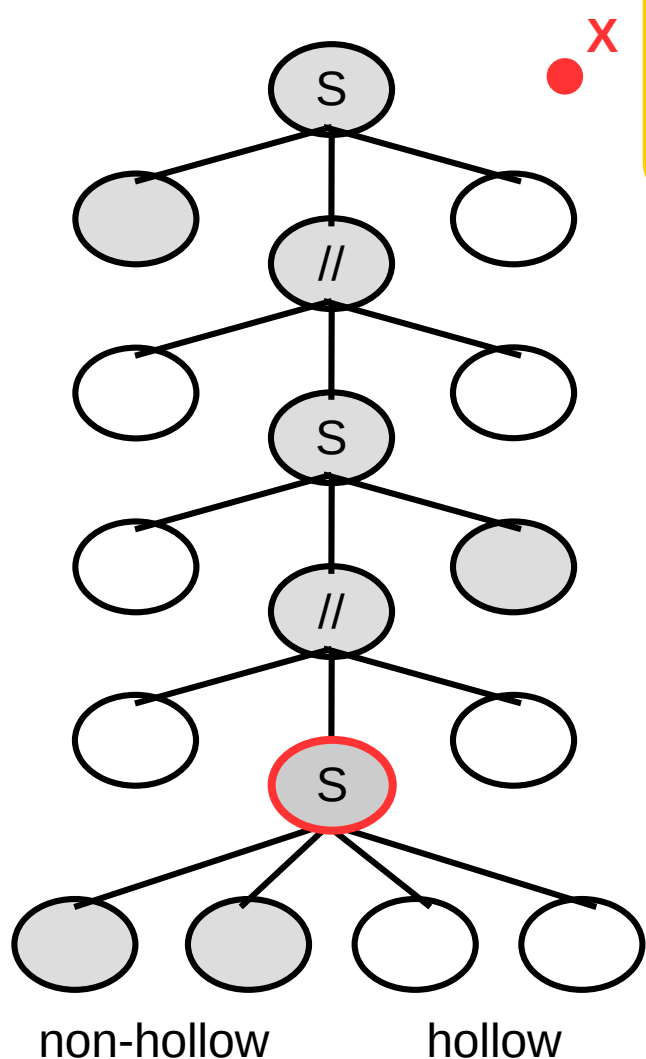
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Definition: u is an **eligible node** iff all parallel strict ancestors of u are such that all their children (but one) are hollow

Completion anchored at u

In our algorithm : $G+x$ is not a cograph



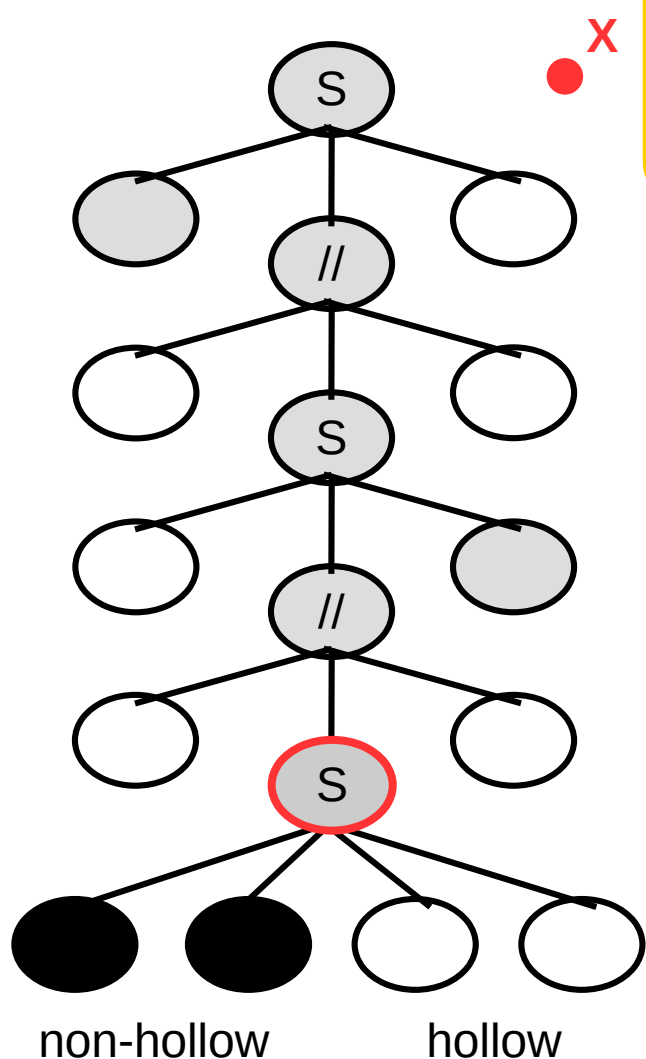
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Proceed as follows :

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Completion anchored at u

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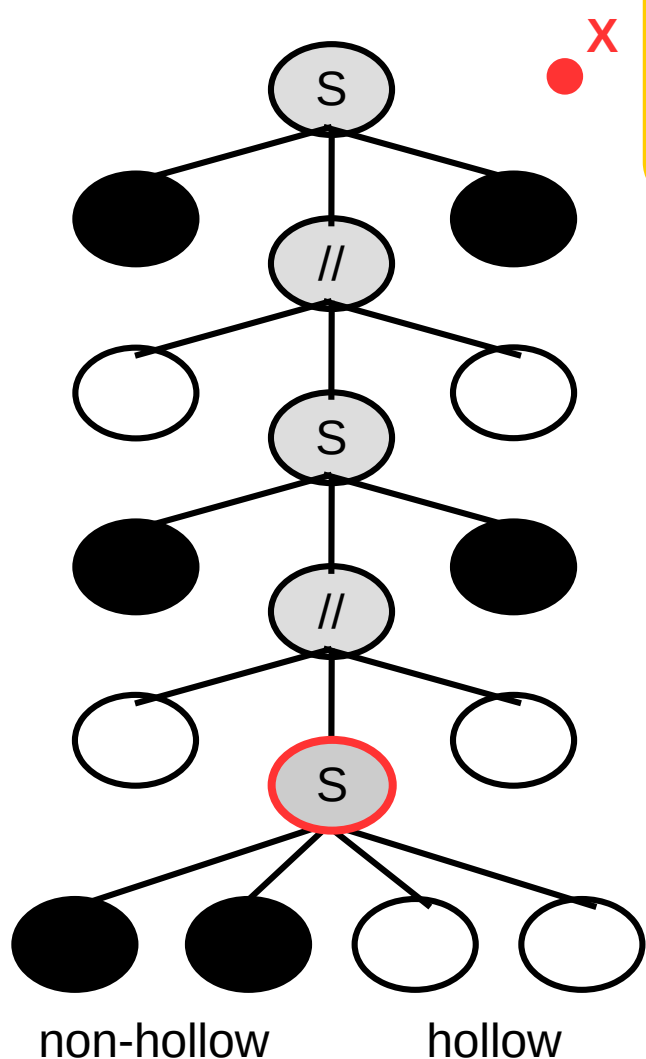
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Proceed as follows :

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Completion anchored at u

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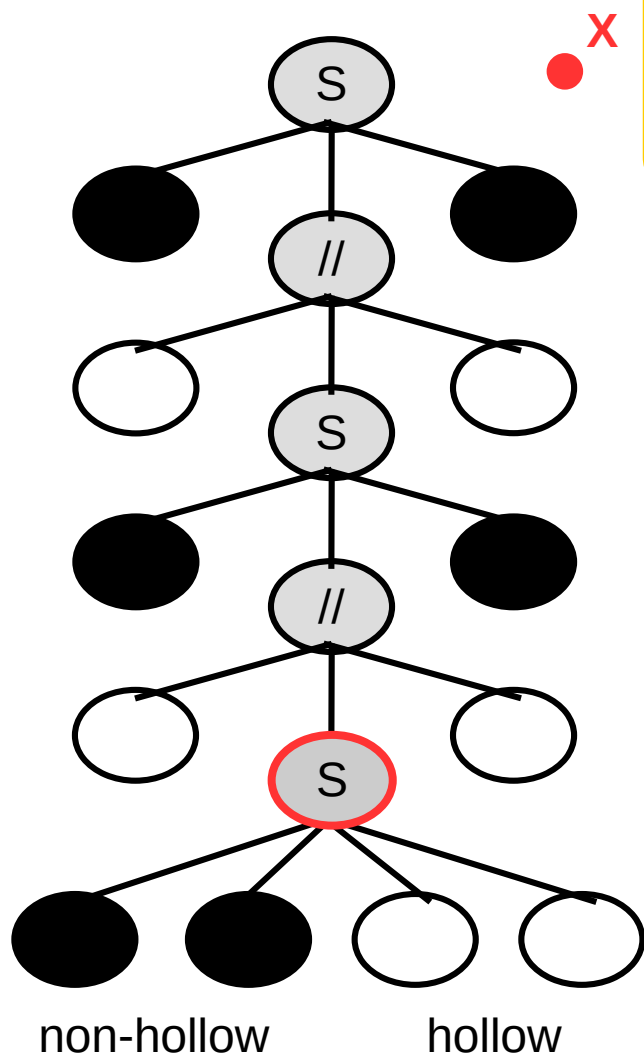
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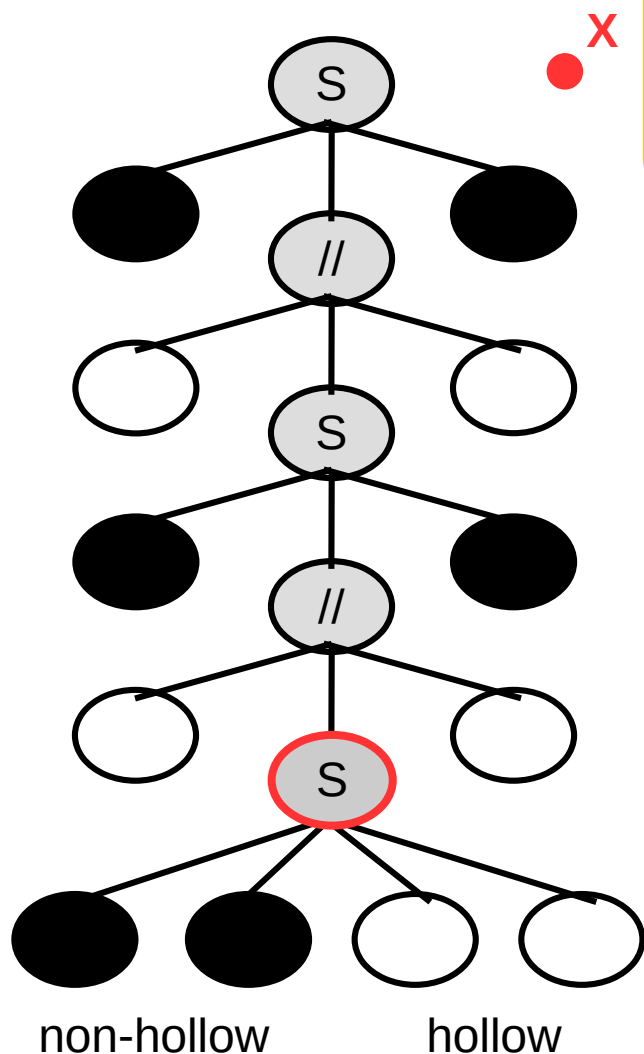
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➔ you obtain a cograph completion of $G+x$
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Completion anchored at u

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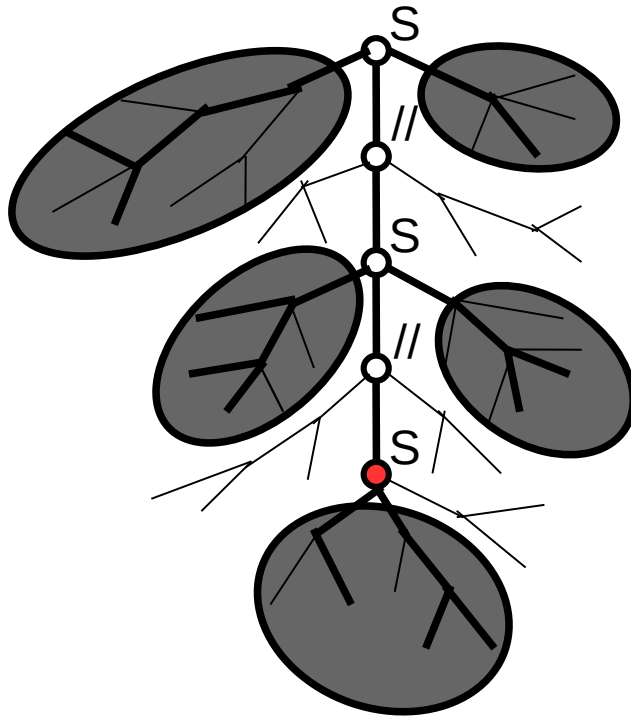
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called the **completion anchored at u**

Question: Is it minimal ?

➔ We have a characterization for this

First algorithm : $O(n+m')$

- Search the tree bottom up from the leaves adjacent to x
- Find the eligible nodes that satisfy the characterization



Note : we search only non-hollow nodes

Complexity : $O(d')$

[LMP 10]

- Choose one u of **minimum cost** and update the data structure by running [CPS 81]'s algorithm.

Complexity : $O(d')$ for one incremental step
 $O(n+m')$ for the whole algorithm

Completion algorithms

Second algorithm: $O(n + m \log^2 n)$

Why is $O(n+m')$ not necessarily optimal?

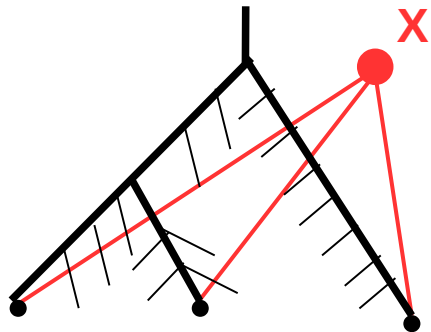
- No reason to use adjacency lists to encode the output
 - there is an $O(n)$ space representation of cographs

Why is $O(n+m')$ not necessarily optimal?

- No reason to use adjacency lists to encode the output
 - ➔ there is an $O(n)$ space representation of cographs
- What is the expected number of edges m' in a cograph completion?
 - If the input G has the vertex-expansion property, then G' has $O(n^2)$ edges
 - Random graphs with fixed average degree, **$O(n)$ edges**, have the expansion property with high probability
 - ➔ **In practice, $O(n+m') \sim O(n^2)$**
 - ➔ **We achieve $O(n+m \log^2 n)$ time**

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 - ➔ In practice, $O(n+m') \sim O(n^2)$
 - ➔ We achieve $O(n+m \log^2 n)$ time
- Where is the room for improvement of the complexity?



A **constant** number of neighbours of x can force to search an $\Omega(n)$ part of the co tree

Second algorithm : $O(n + m \log^2 n)$

- Note: we abandon the minimum incremental → *only minimal*
- we use a dynamic data-structure for *lowest ancestor queries* [Sleator, Tarjan 1983]
 - In $O(\log n)$ time: $w = \text{lca}(u, v)$ and w_u the child of w that is an ancestor of u
 - Update the structure in $O(\log n)$ time under elementary tree modifications
- we use *ordered lists* [Dietz, Sleator 1987]
 - In $O(1)$ time: order between two elements in the list
 - Update the structure in $O(1)$ time under deletion and insertion of an element

Second algorithm : $O(n + m \log^2 n)$

Our goal : determine the *lowest* eligible, non-hollow and non-forced nodes

→ **minimal completion**

Second algorithm : $O(n + m \log^2 n)$

Our goal : determine the *lowest* eligible, non-hollow and non-forced nodes

→ minimal completion

■ Lowest eligible nodes

→ highest parallel nodes with ≥ 2 non-hollow children

Second algorithm : $O(n + m \log^2 n)$

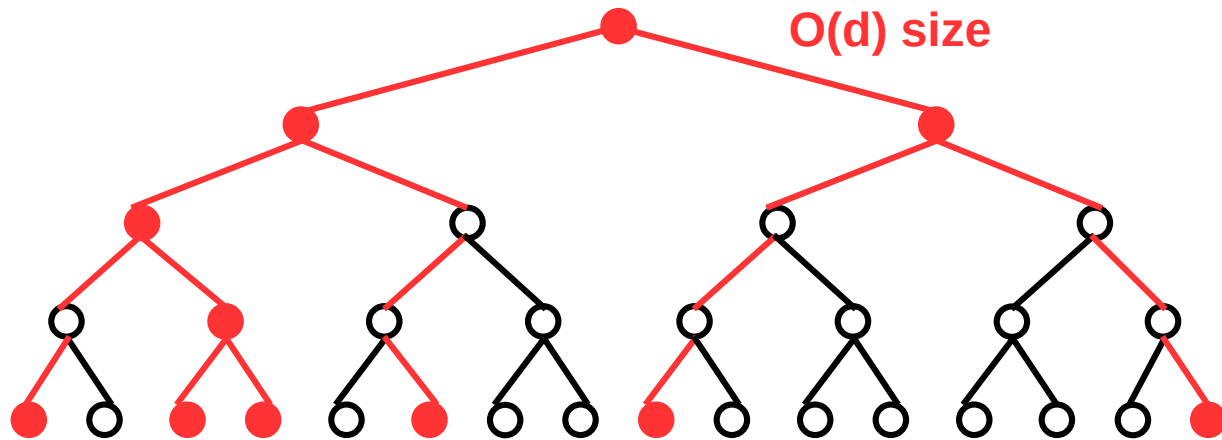
Our goal : determine the *lowest* eligible, non-hollow and non-forced nodes

→ minimal completion

■ Lowest eligible nodes

→ highest parallel nodes with ≥ 2 non-hollow children

- build T' : the subtree of lowest common ancestors of neighbors of x
- Keep the highest parallel nodes in T'



Second algorithm : $O(n + m \log^2 n)$

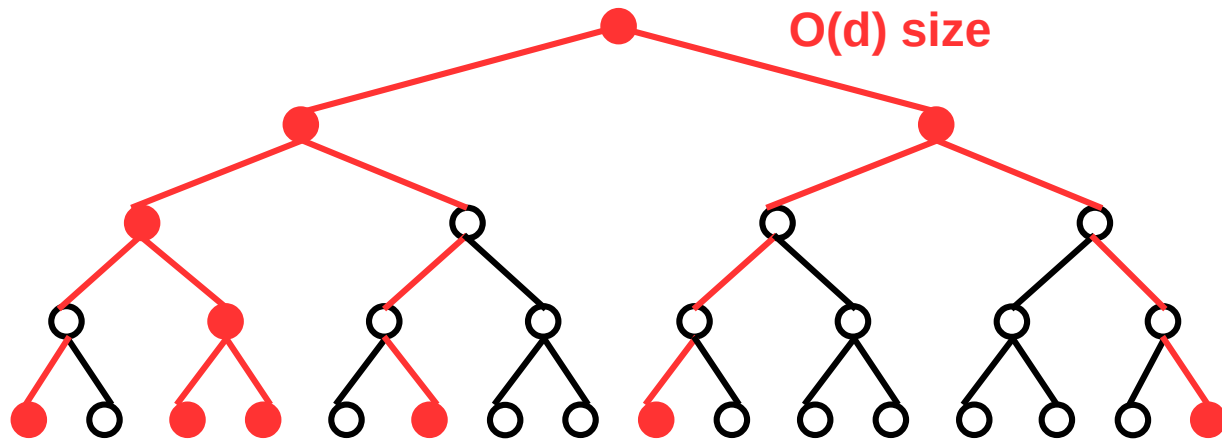
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■ Lowest eligible nodes

→ highest parallel nodes with ≥ 2 non-hollow children

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1) sort neighbours of x from left to right : $O(d \log^2 n)$ time

Second algorithm : $O(n + m \log^2 n)$

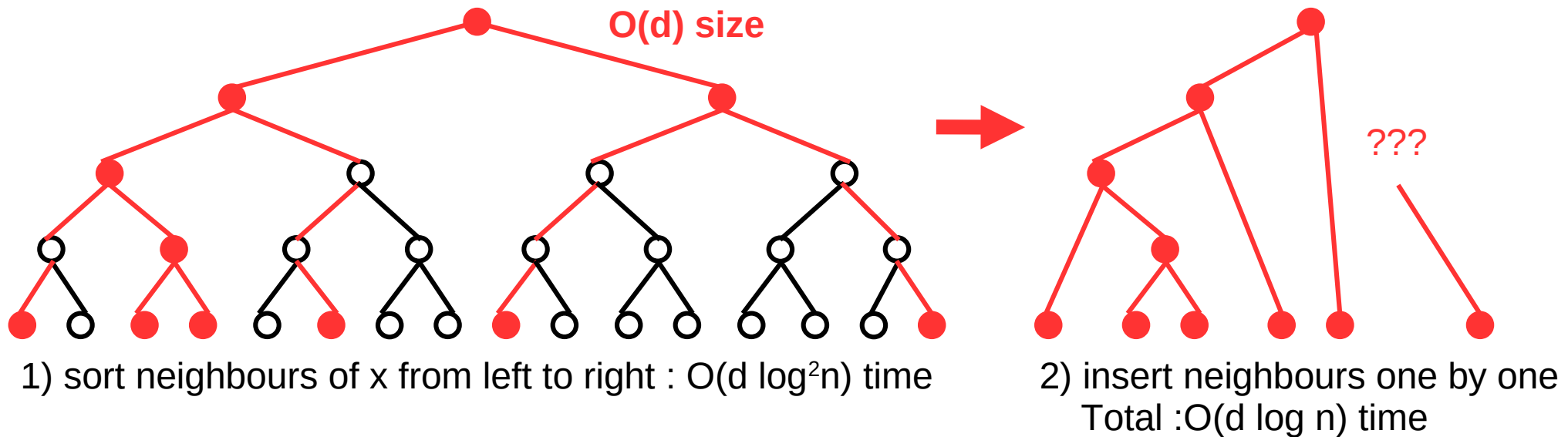
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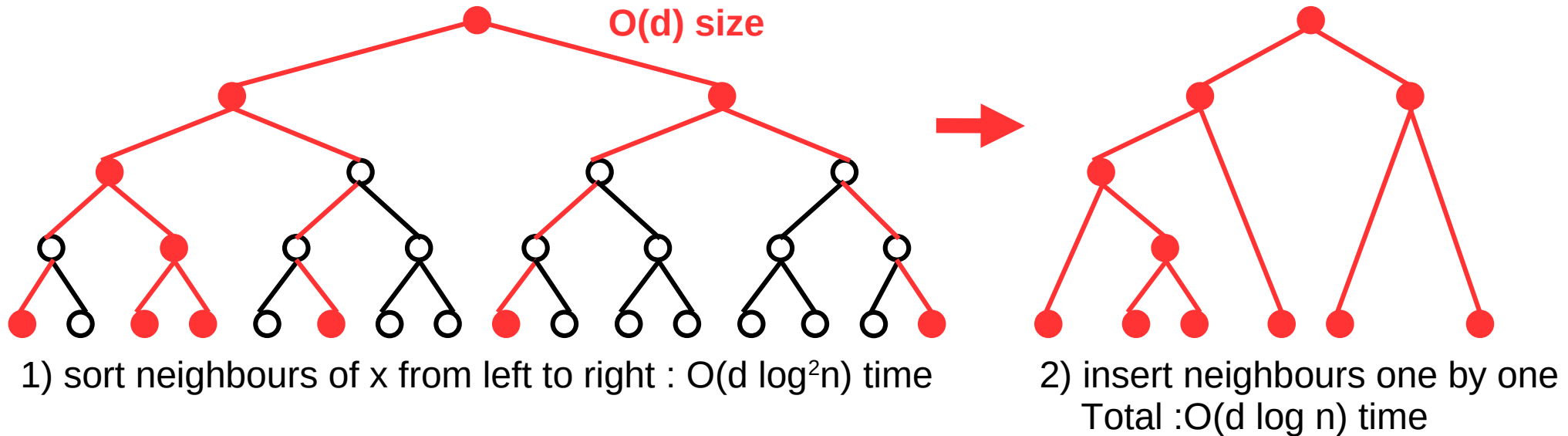
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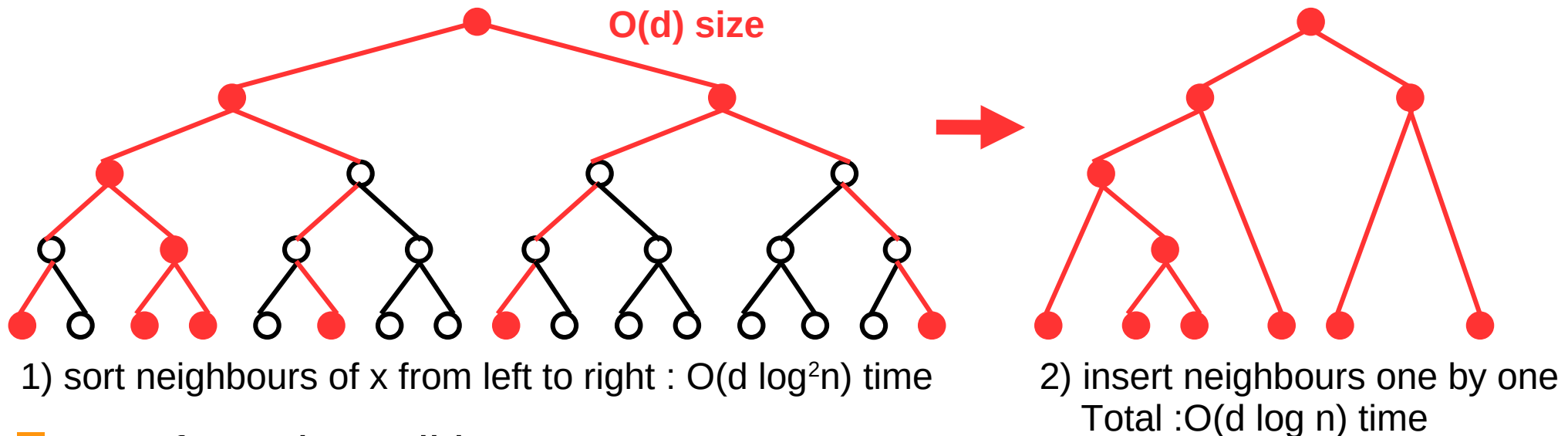
Our goal : determine the *lowest* eligible, non-hollow and non-forced nodes

→ **minimal completion**

■ Lowest eligible nodes

→ **highest parallel nodes with ≥ 2 non-hollow children**

- build T' : the subtree of lowest common ancestors of neighbors of x
- Keep the highest parallel nodes in T'



■ Non-forced condition

- Find the lowest non-forced node above each node of W (grand-parent)

Complexity : $O(d \log^2 n)$ for one incremental step
 $O(n + m \log^2 n)$ for the whole algorithm

Editing algorithm

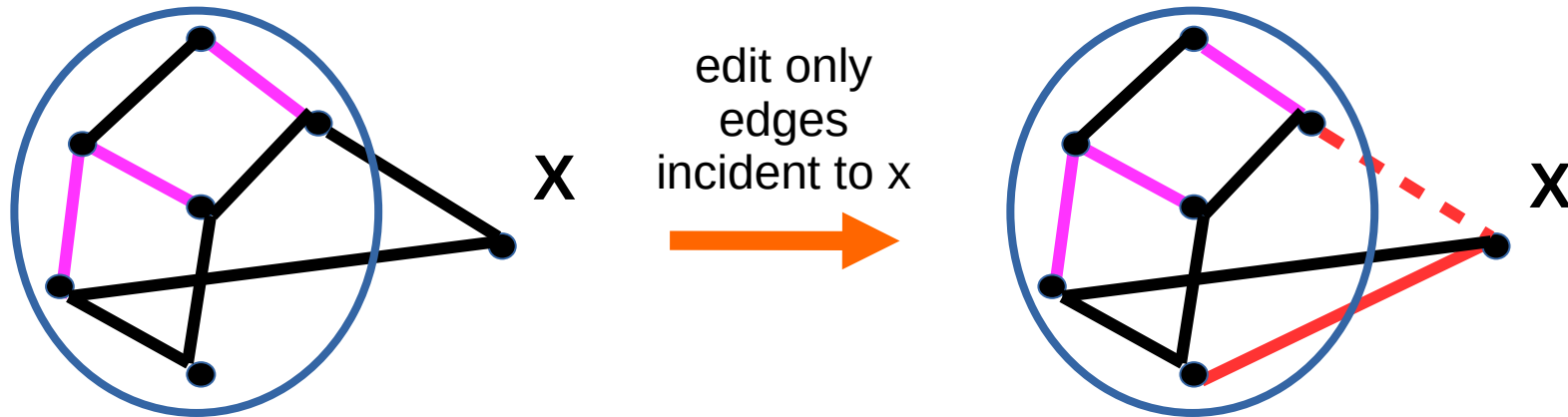
$O(n + m)$ time

Algorithm for cograph editing

- Editing: use both additions and deletions of edges
- Minimal for inclusion
- **Linear time: $O(n+m)$**
- Additional feature: **minimum editing at each incremental step**
- **number of edits returned is $\leq m$**

The local incremental approach

- Vertices are processed one by one
- Only edges *incident to x* are modified



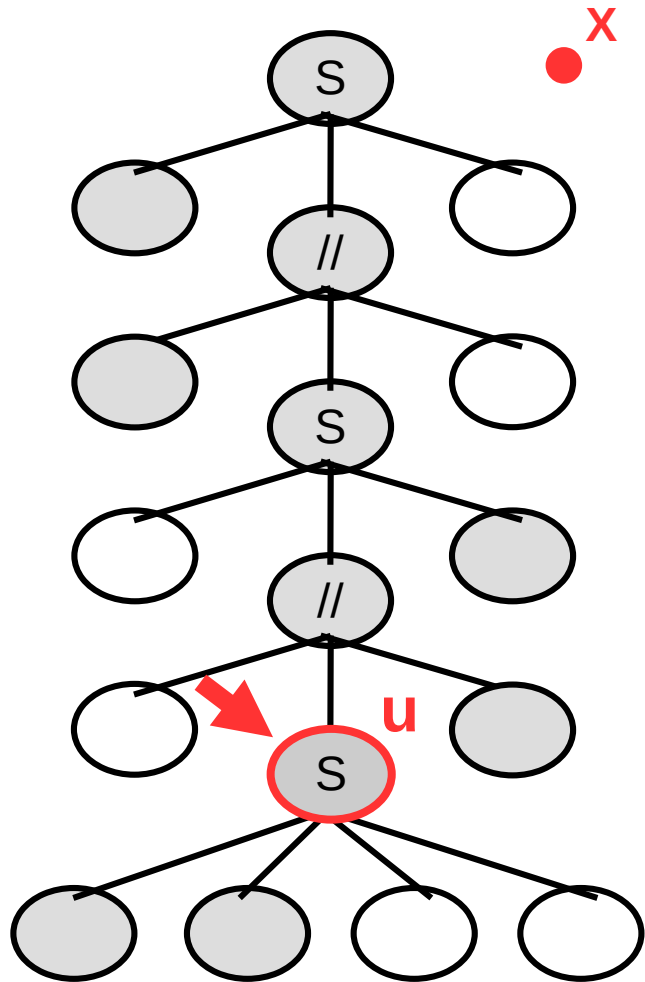
Always possible when:

- The class is *hereditary*
- Contains no maximal element for induced subgraph relationship

➔ Our goal : $O(d)$ time complexity at each incremental step

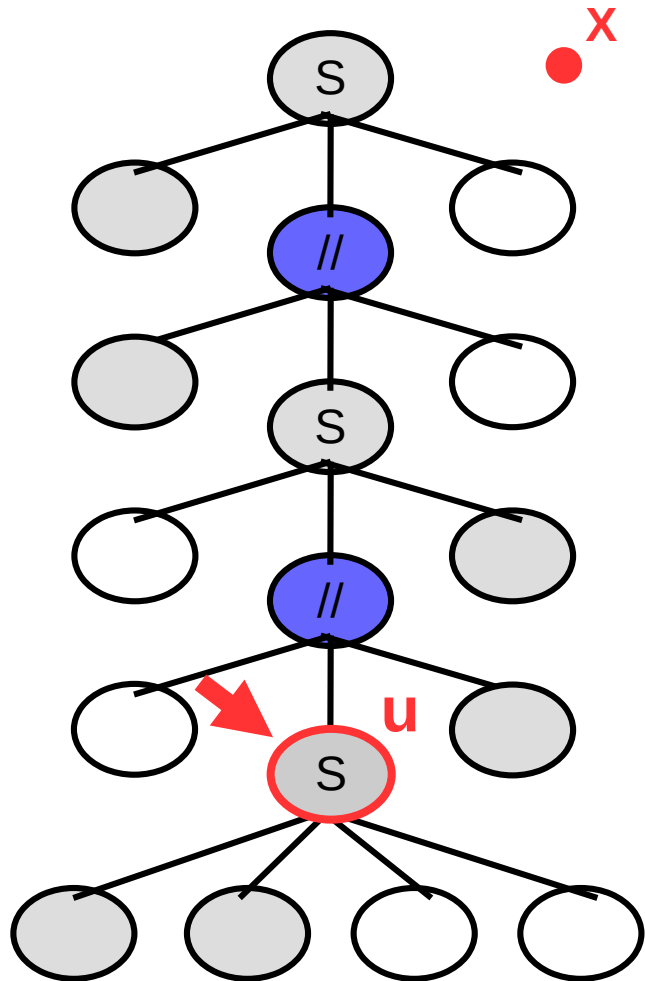
Editing anchored at u

In our algorithm : $G+x$ is not a cograph



Editing anchored at u

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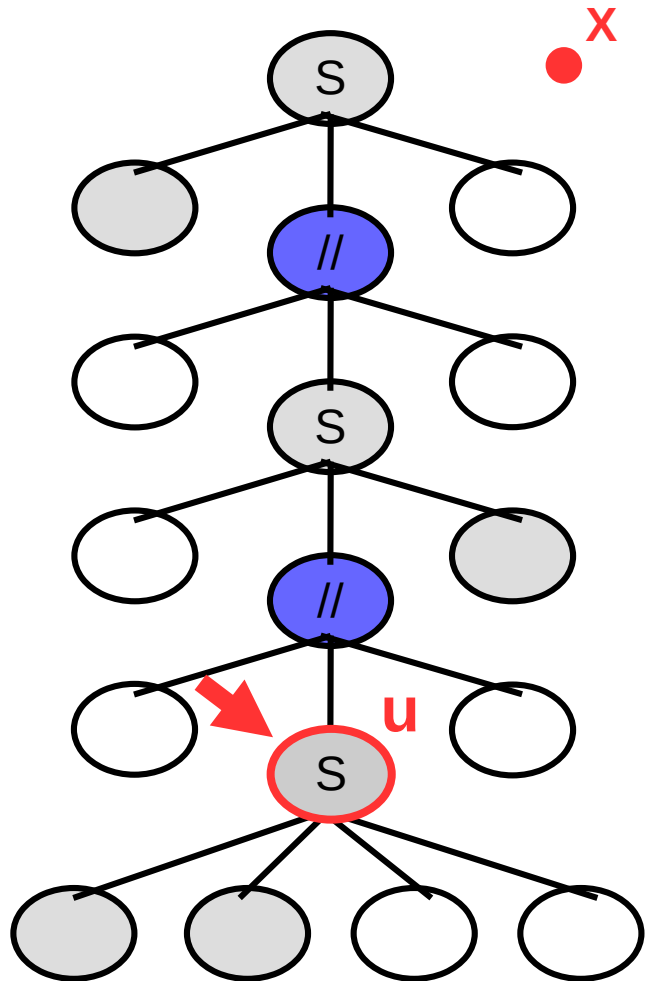


Proceed as follows :

- 1) for each **parallel ancestor** of u , make all its children (but one) **hollow**

Editing anchored at u

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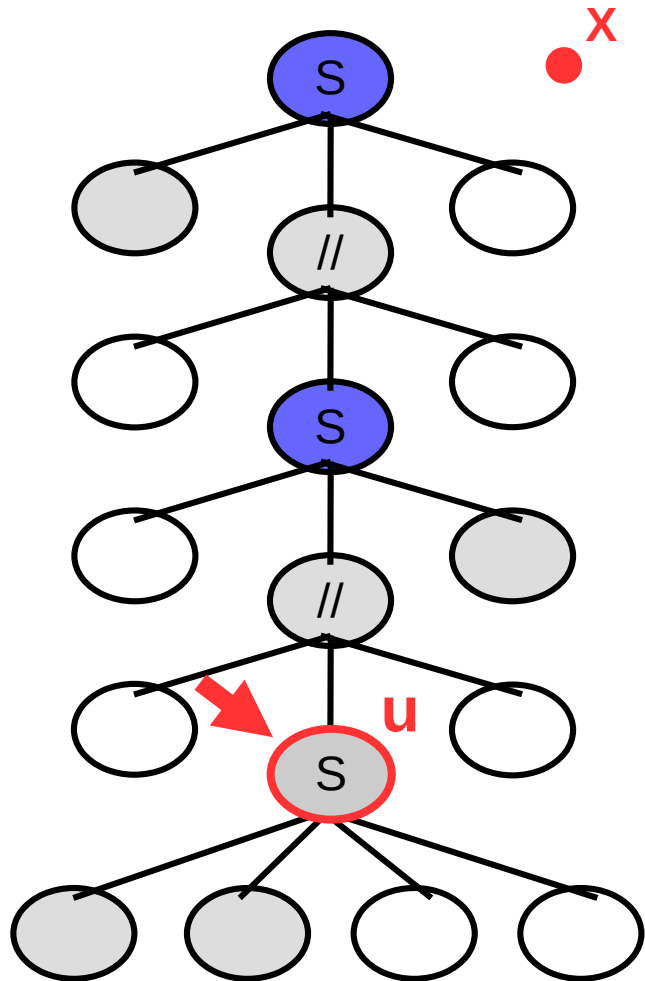


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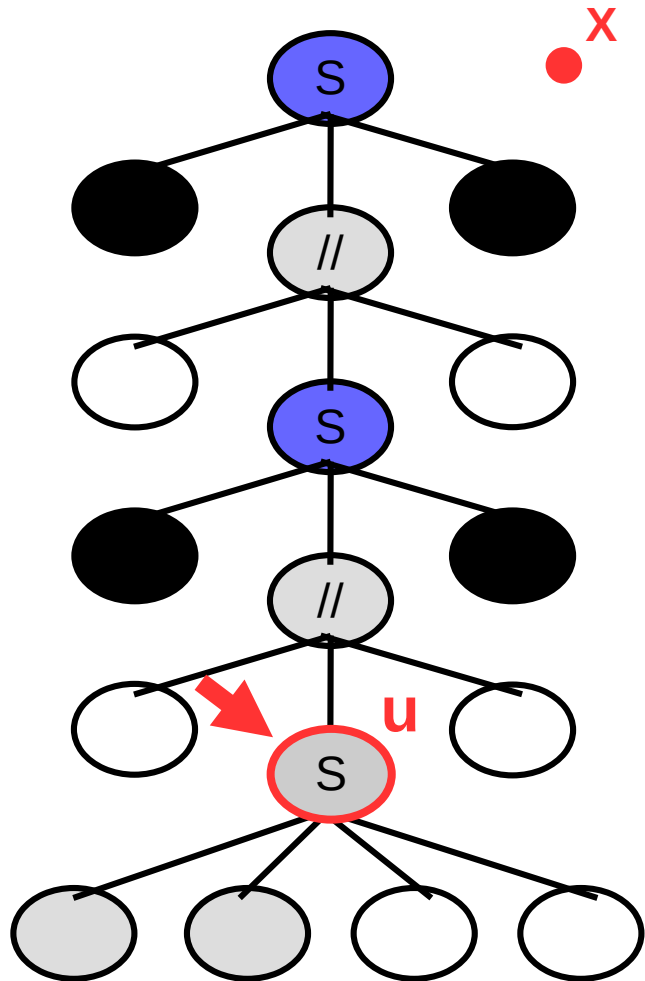


Proceed as follows :

- 1) for each **parallel ancestor** of u , make all its children (but one) **hollow**
- 2) for each **series ancestor** of u , make all its children (but one) **full**

Editing anchored at u

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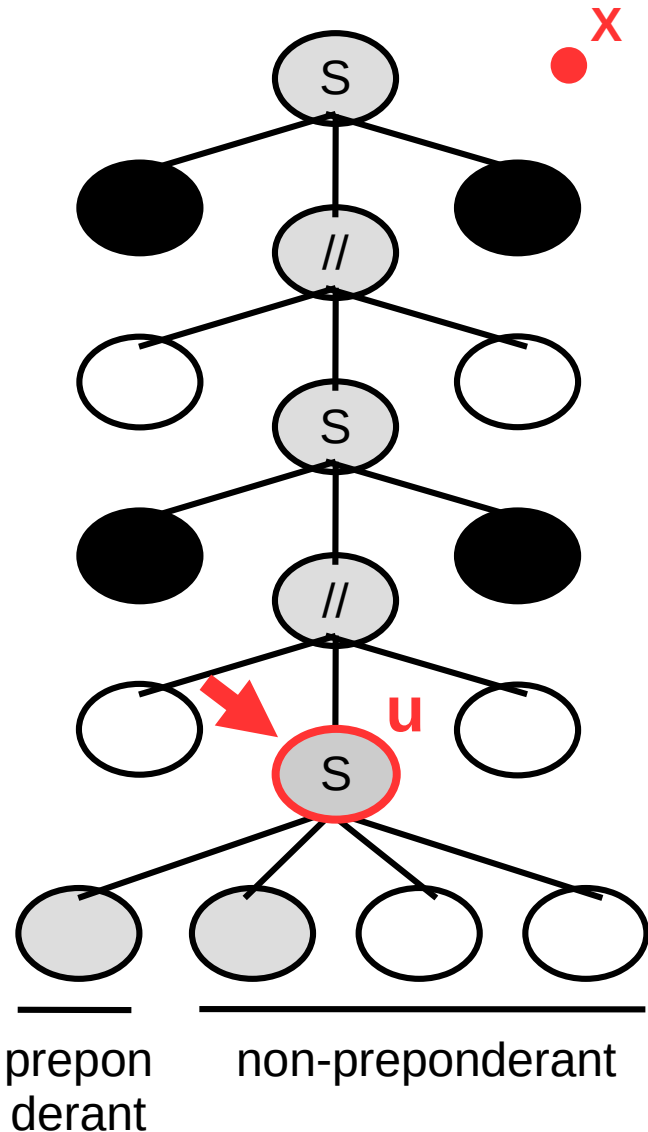


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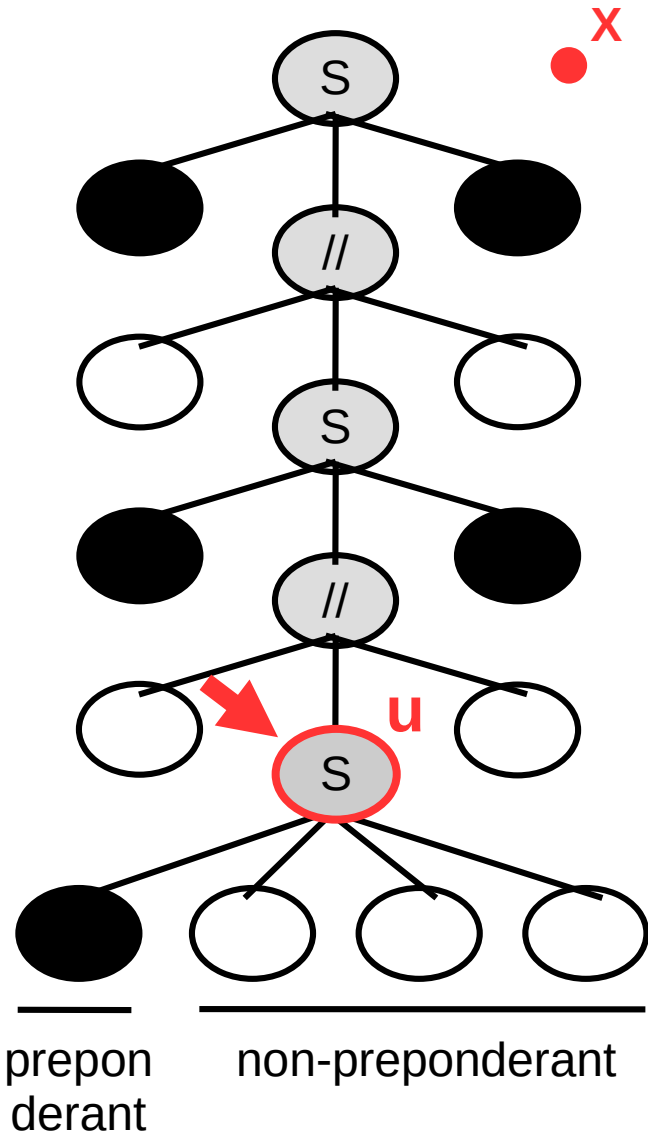


Proceed as follows :

- 1) for each **parallel ancestor** of u , make all its children (but one) **hollow**
- 2) for each **series ancestor** of u , make all its children (but one) **full**
- 3) make the preponderant children of u become **full** and make the non-preponderant ones **hollow**

Editing anchored at u

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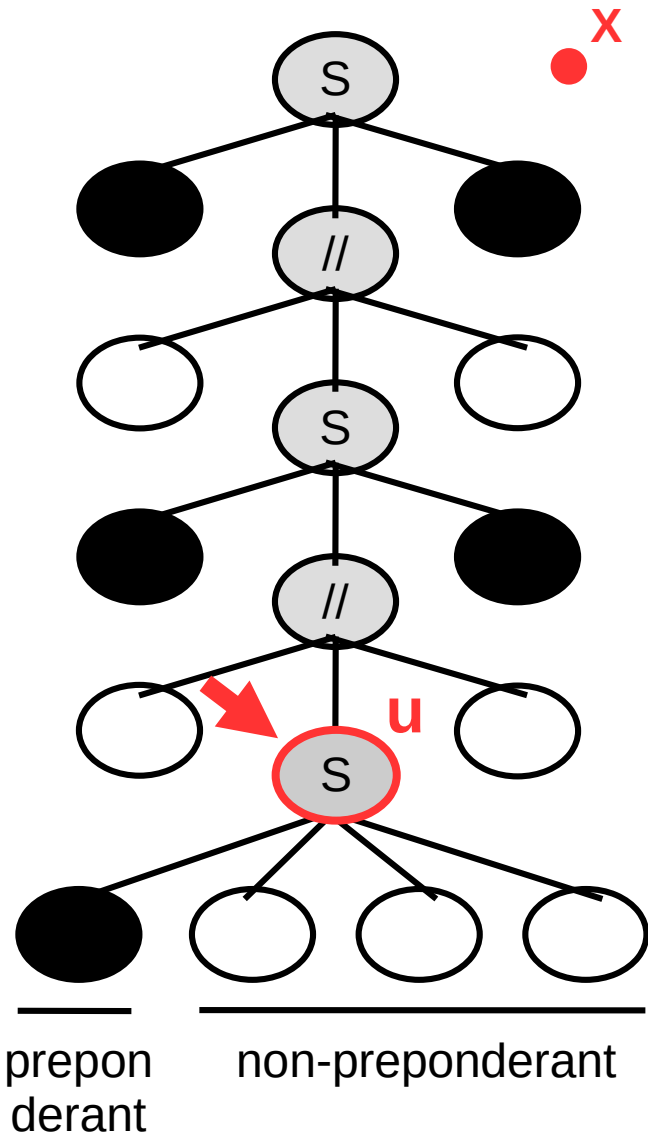
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➔ you obtain a cograph editing of $G+x$
called the **editing anchored at u**

Editing anchored at u

In our algorithm : $G+x$ is not a cograph



Proceed as follows :

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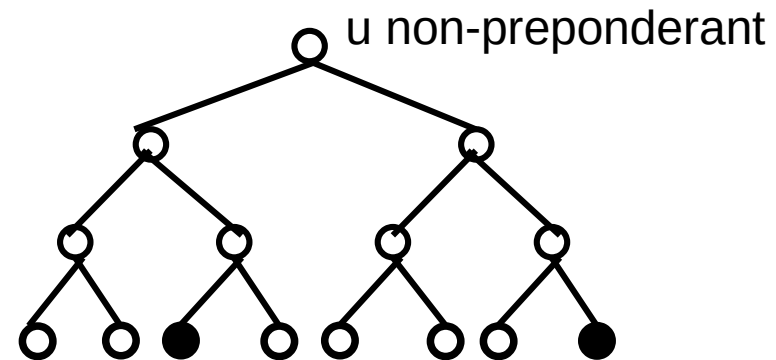
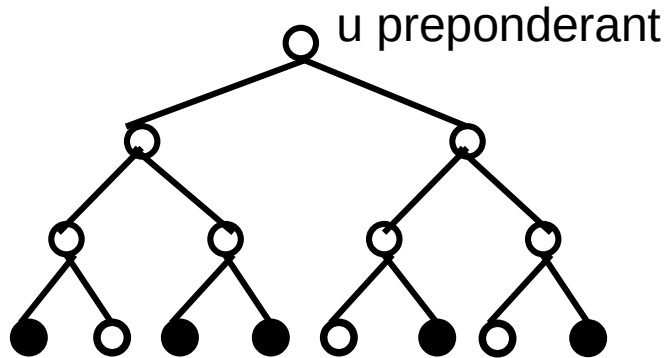
➔ you obtain a cograph editing of $G+x$
called the **editing anchored at u**

Question: Is it minimal? minimum ?

➔ $O(n)$ time algorithm trying all possible nodes of the cotree

Maximal preponderant nodes

Def.: u is preponderant iff the subtree of u contains more neighbours of x than non-neighbours of x



Def.: u is maximal preponderant iff u is preponderant and no ancestor of u is.

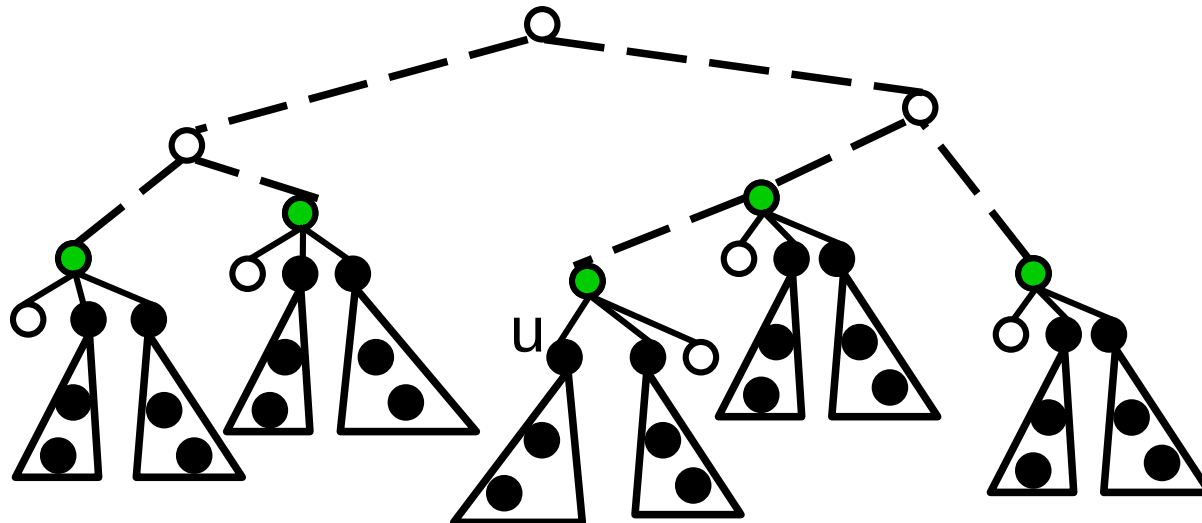
Cor. [CPS81]: the insertion node of a minimum editing has a preponderant child

➔ The insertion node is either in the subtree of some maximal preponderant node or is the parent of some maximal preponderant node

Only $O(d)$ candidates for the insertion node

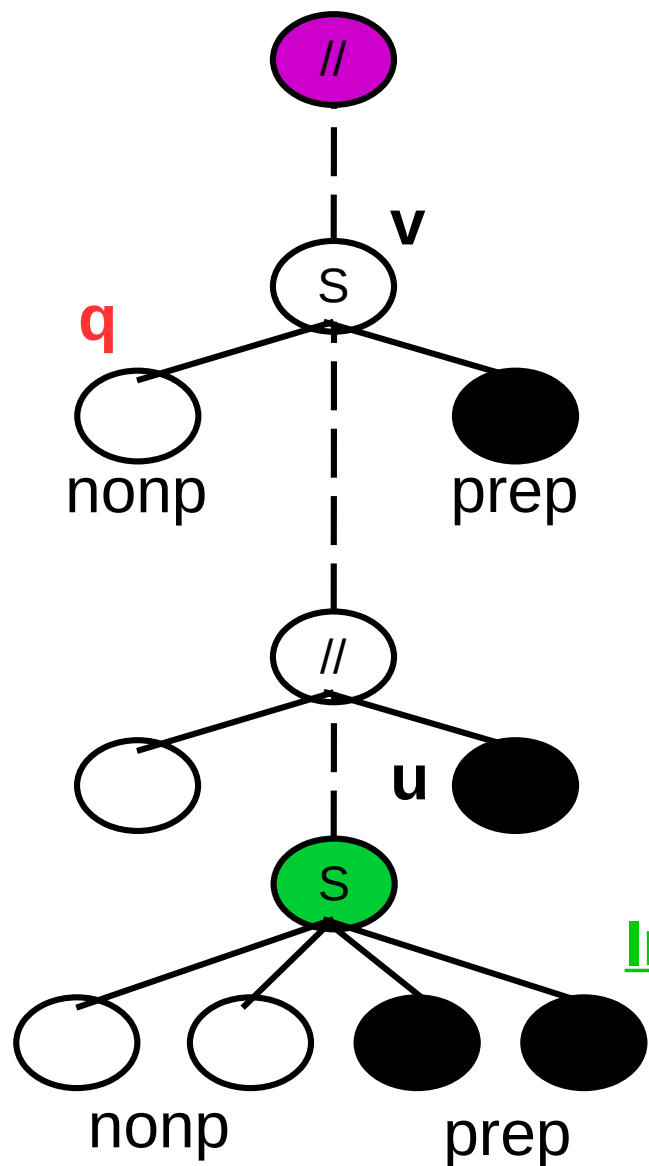
Outline of the algorithm

- 1) compute all maximal *preponderant* nodes (and their parents)
- 2) for each maximal preponderant node u , determine the minimum editing anchored in its subtree or at its parent
 - ➔ $O(n)$ algo applied on a subtree where $n=O(d)$
- 3) keep the minimum editing among all the editings found for each maximal preponderant node u : need to compute ***cost-above(u)***



Principle of the bottom-up search

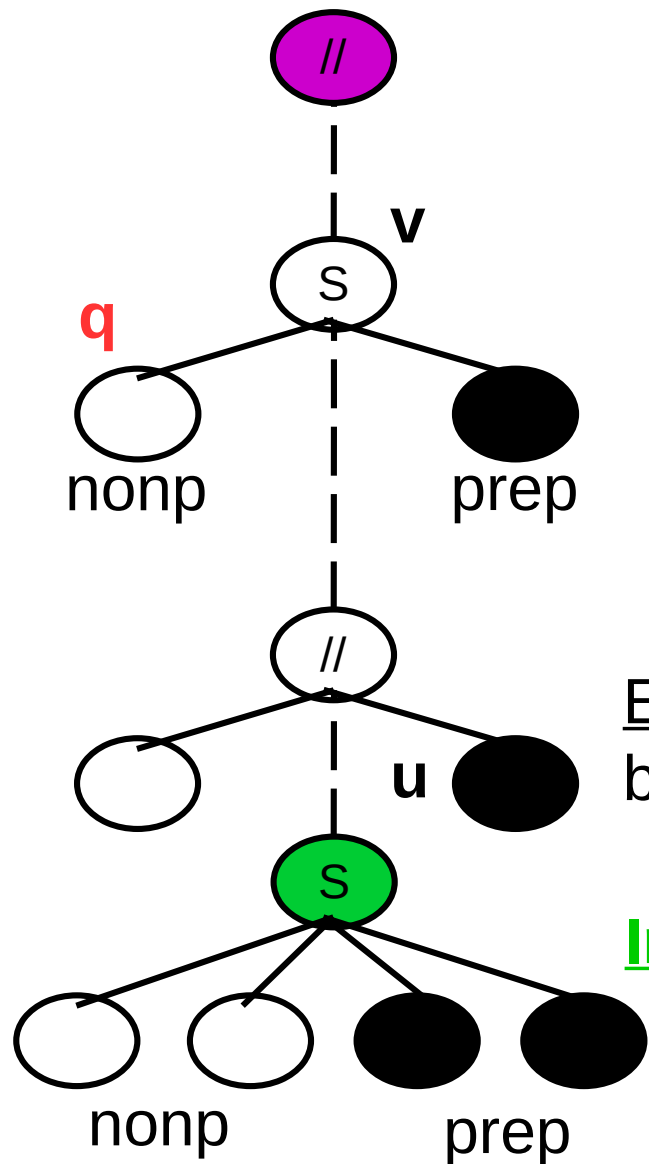
Obs.: we need the cost of the editing anchored at u only if it is less than the cost of the delete-all editing



Initialisation :
$$\text{bud}(u) = \frac{B_{\text{prep}}(u) - W_{\text{prep}}(u)}{= \text{exc}(u)}$$

Principle of the bottom-up search

Obs.: we need the cost of the editing anchored at u only if it is less than the cost of the delete-all editing

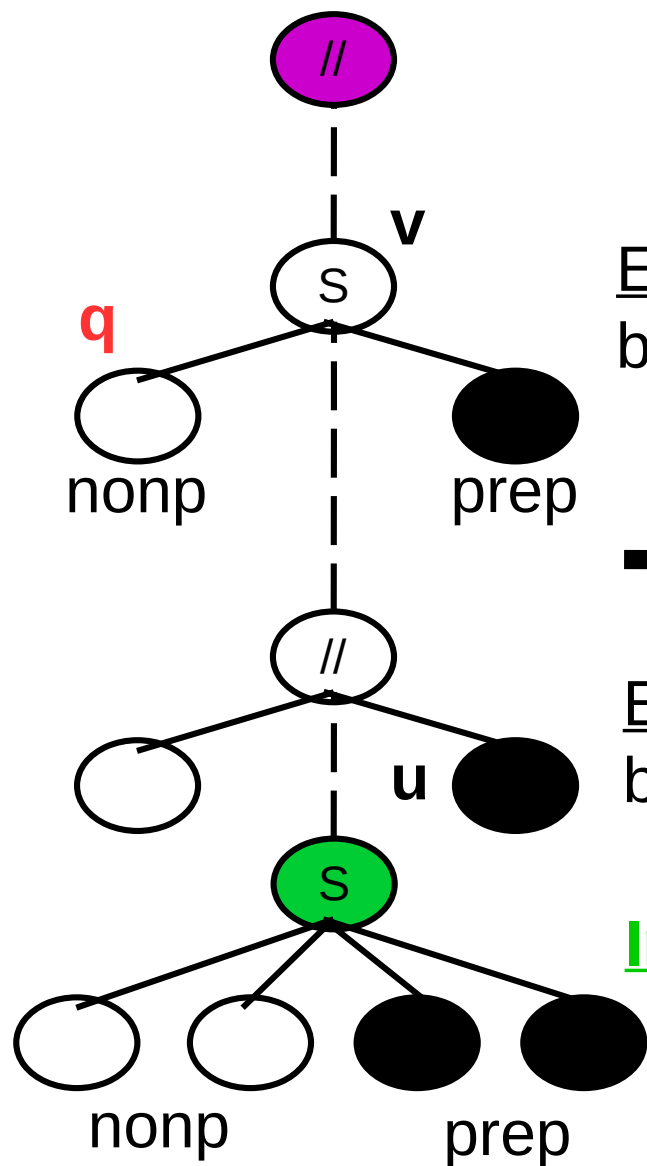


Encounter a parallel node:
 $\text{bud}(u)$ unchanged

Initialisation : $\text{bud}(u) = \frac{B_{\text{prep}}(u) - W_{\text{prep}}(u)}{= \text{exc}(u)}$

Principle of the bottom-up search

Obs.: we need the cost of the editing anchored at u only if it is less than the cost of the delete-all editing



Encounter a series node v :

$$\text{bud}(u) \leftarrow \frac{\text{bud}(u) + B_{\text{prep}}(v) - W_{\text{prep}}(v) + B_{\text{nonp}}(v) - W_{\text{nonp}}(v)}{1} = \text{bud}$$

➔ Routine SearchTree(**q**, **bud**)

Encounter a parallel node:

$\text{bud}(u)$ unchanged

Initialisation : $\text{bud}(u) = \frac{B_{\text{prep}}(u) - W_{\text{prep}}(u)}{1} = \text{exc}(u)$

Principle of the bottom-up search

Obs.: we need the cost of the editing anchored at u only if it is less than the cost of the delete-all editing

Stop when either the budget becomes negative or when the search reaches the root with non-negative budget \rightarrow deduce $\text{cost-above}(u)$

Encounter a series node v :

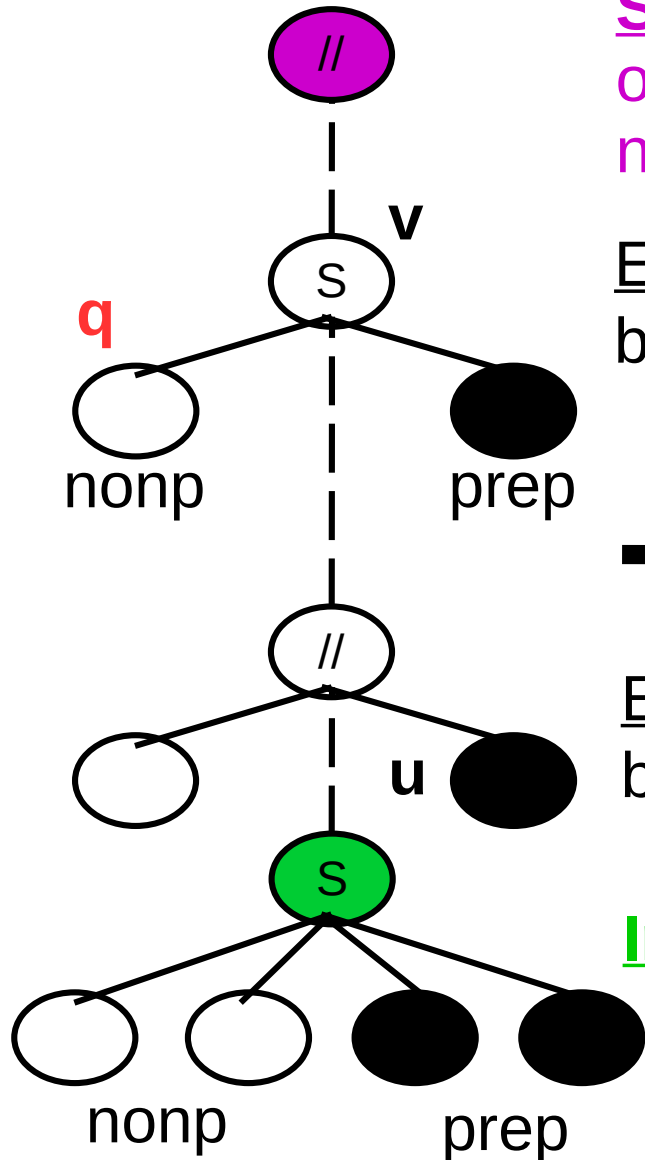
$$\text{bud}(u) \leftarrow \frac{\text{bud}(u) + B_{\text{prep}}(v) - W_{\text{prep}}(v) + B_{\text{nonp}}(v) - W_{\text{nonp}}(v)}{=} = \text{bud}$$

\rightarrow Routine $\text{SearchTree}(\mathbf{q}, \mathbf{bud})$

Encounter a parallel node:

$\text{bud}(u)$ unchanged

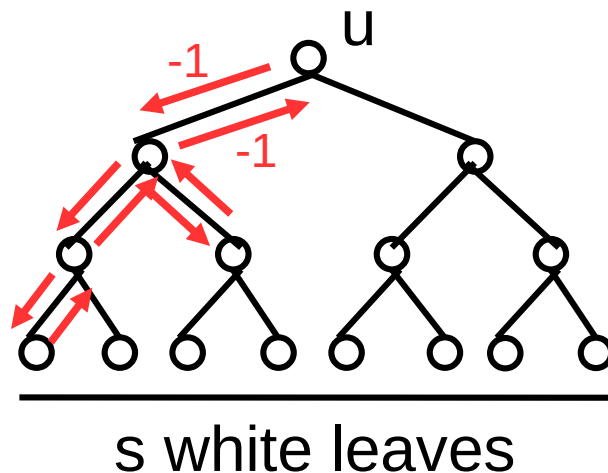
Initialisation : $\text{bud}(u) = \frac{B_{\text{prep}}(u) - W_{\text{prep}}(u)}{=} = \text{exc}(u)$



Routine SearchTree(u,s)

- Makes a **DFS limited by a *tll*** and counts the difference between black and white leaves in *cpt*
 - Initially, $tll \leftarrow 2+5s$ and $cpt \leftarrow s$
 - *tll* is decreased when an edge is traversed
 - DFS stops when $tll=-1$
- Main property:
 $W(u)-B(u) \leq s$ iff Search-tree(*u*, *s*) searches the entire subtree of *u* and ends with a value $cpt \geq 0$.
Complexity : **$O(\min\{s, W(u)-B(u)\})$**

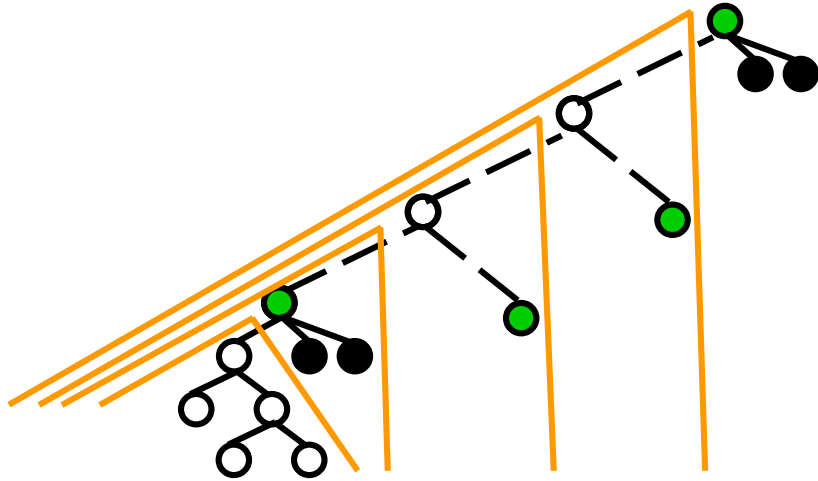
why $tll \leftarrow 2+5s$?



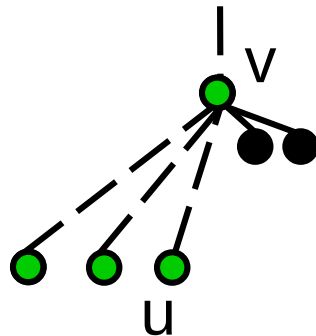
4s-2 edge traversals

Two threats to the complexity

- Searching repeatedly the same part of the tree with the same budget



- Using repeatedly the same budget in the bottom-up search



$$\text{bud}(u) \leftarrow \text{bud}(u) + \mathbf{B}_{\text{prep}}(\mathbf{v}) - \mathbf{W}_{\text{prep}}(\mathbf{v}) + \mathbf{B}_{\text{nonp}}(\mathbf{v}) - \mathbf{W}_{\text{nonp}}(\mathbf{v})$$

Some open algorithmic questions

Inclusion-minimal cograph editing in linear time

- minimum at each incremental step
- at most m edits at the end

- Showing that minimal cograph completion is ***not*** solvable in linear time
 $O(n+m\log^2n)$ from [Crespelle, Lokshtanov, Phan, Thierry 2020]
- Inclusion-minimal editing for other graph classes, in linear time?

**Complex networks as
almost cographs?**

Cograph edition of real-world graphs

35 real-world
graphs

+

8 random
graphs

Context	Network	n	m	d ^o	%mod
WWW	in-2004	1 148 875	12 281 937	21.4	12 %
WWW	cnr-2000	227 058	2 187 201	19.3	19 %
PROTEIN	reactome	5 973	145 778	48.8	22 %
SOFTWARE	jdk	6 434	53 658	16.7	29 %
SOFTWARE	jung-j	6 120	50 290	16.4	29 %
WWW	eu-2005	835 044	15 718 784	37.7	29 %
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SPECIES	foodweb	183	2 434	26.6	43 %
CO-AUTHOR	dblp	317 080	1 049 866	6.6	45 %
WORD-REL.	wordnet	145 145	656 230	9.0	48 %
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CO-OCCUR	bible-names	1 707	9 059	10.6	67 %
PROTEIN	figeys	2 217	6 418	5.8	67 %
CITATION-SCI.	cora	23 166	89 157	7.7	68 %
SOCIAL	youtube	1 134 890	2 987 624	5.3	69 %
CO-ACTOR	actor-col.	374 511	15 014 839	80.2	71 %
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RANDOM	ER-Gnm_1M-8	999 684	3 999 999	8.0	84 %
RANDOM	ER-Gnm_1M-10	999 952	5 000 000	10.0	87 %
RANDOM	ER-Gnm_1M-15	1 000 000	7 500 000	15.0	91 %
SOCIAL	orkut	3 072 441	117 185 083	76.3	91 %
RANDOM	ER-Gnm_1M-20	1 000 000	10 000 000	20.0	93 %
WORD-REL.	Thesaurus	23 132	297 094	25.7	93 %

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RESULTS

- Some networks are very close from cographs

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RESULTS

- Some networks are very close from cographs
- Random graphs are never

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35 real-world
graphs

+

8 random
graphs

Context	Network	n	m	d ^o	%mod
WWW	in-2004	1 148 875	12 281 937	21.4	12 %
WWW	cnr-2000	227 058	2 187 201	19.3	19 %
PROTEIN	reactome	5 973	145 778	48.8	22 %
SOFTWARE	jdk	6 434	53 658	16.7	29 %
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SPECIES	foodweb	183	2 434	26.6	43 %
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RANDOM	ER-Gnm_1M-3	940 987	1 494 643	3.2	63 %
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CO-OCCUR	bible-names	1 707	9 059	10.6	67 %
PROTEIN	figeys	2 217	6 418	5.8	67 %
CITATION-SCI.	cora	23 166	89 157	7.7	68 %
SOCIAL	youtube	1 134 890	2 987 624	5.3	69 %
CO-ACTOR	actor-col.	374 511	15 014 839	80.2	71 %
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RANDOM	ER-Gnm_1M-8	999 684	3 999 999	8.0	84 %
RANDOM	ER-Gnm_1M-10	999 952	5 000 000	10.0	87 %
RANDOM	ER-Gnm_1M-15	1 000 000	7 500 000	15.0	91 %
SOCIAL	orkut	3 072 441	117 185 083	76.3	91 %
RANDOM	ER-Gnm_1M-20	1 000 000	10 000 000	20.0	93 %
WORD-REL.	Thesaurus	23 132	297 094	25.7	93 %

RESULTS

- Some networks are very close from cographs
- Random graphs are never
- A wide range of proximity : 12% to 93%

Cograph edition of real-world graphs

35 real-world
graphs
+
8 random
graphs

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Cograph edition of real-world graphs

Close to cographs

- WWW
- software



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Not close not far



 internet
 road

- The proximity with cographs highly depends on the real-world context

Cograph edition of real-world graphs

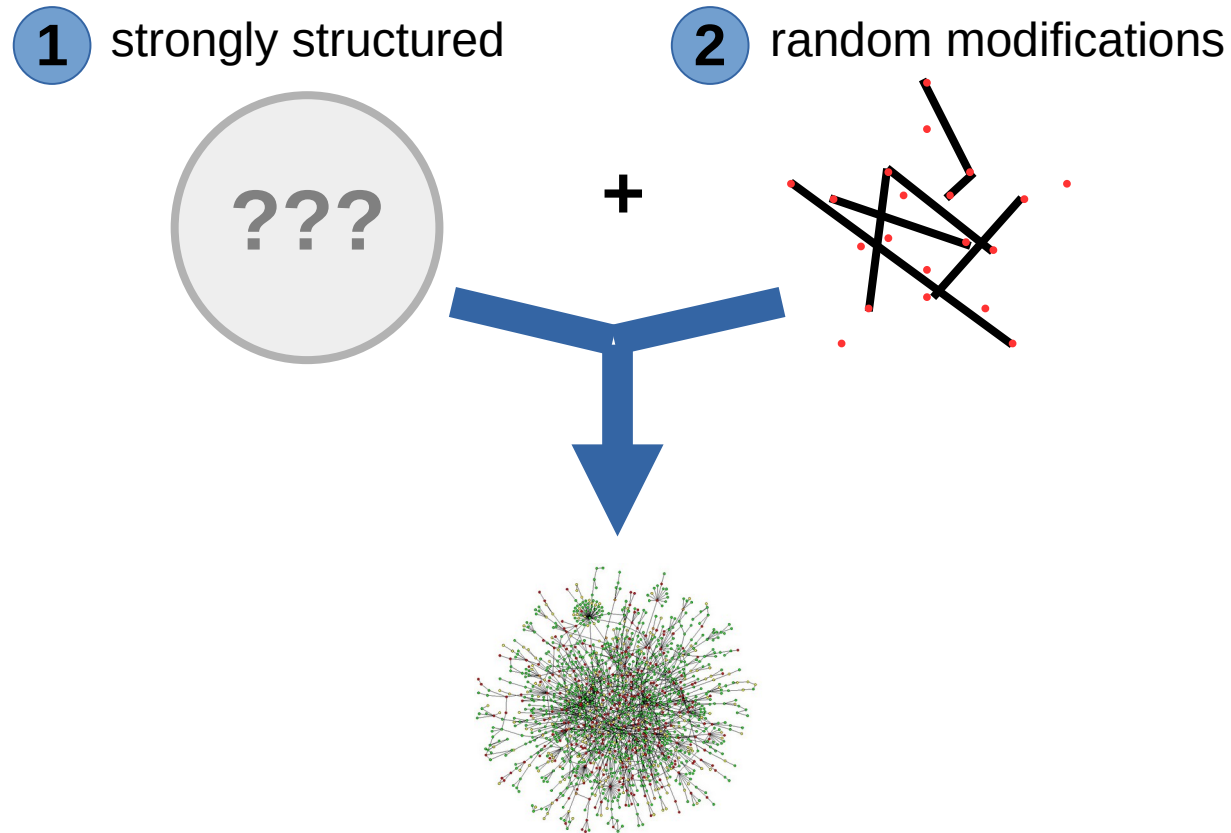
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Far from cographs

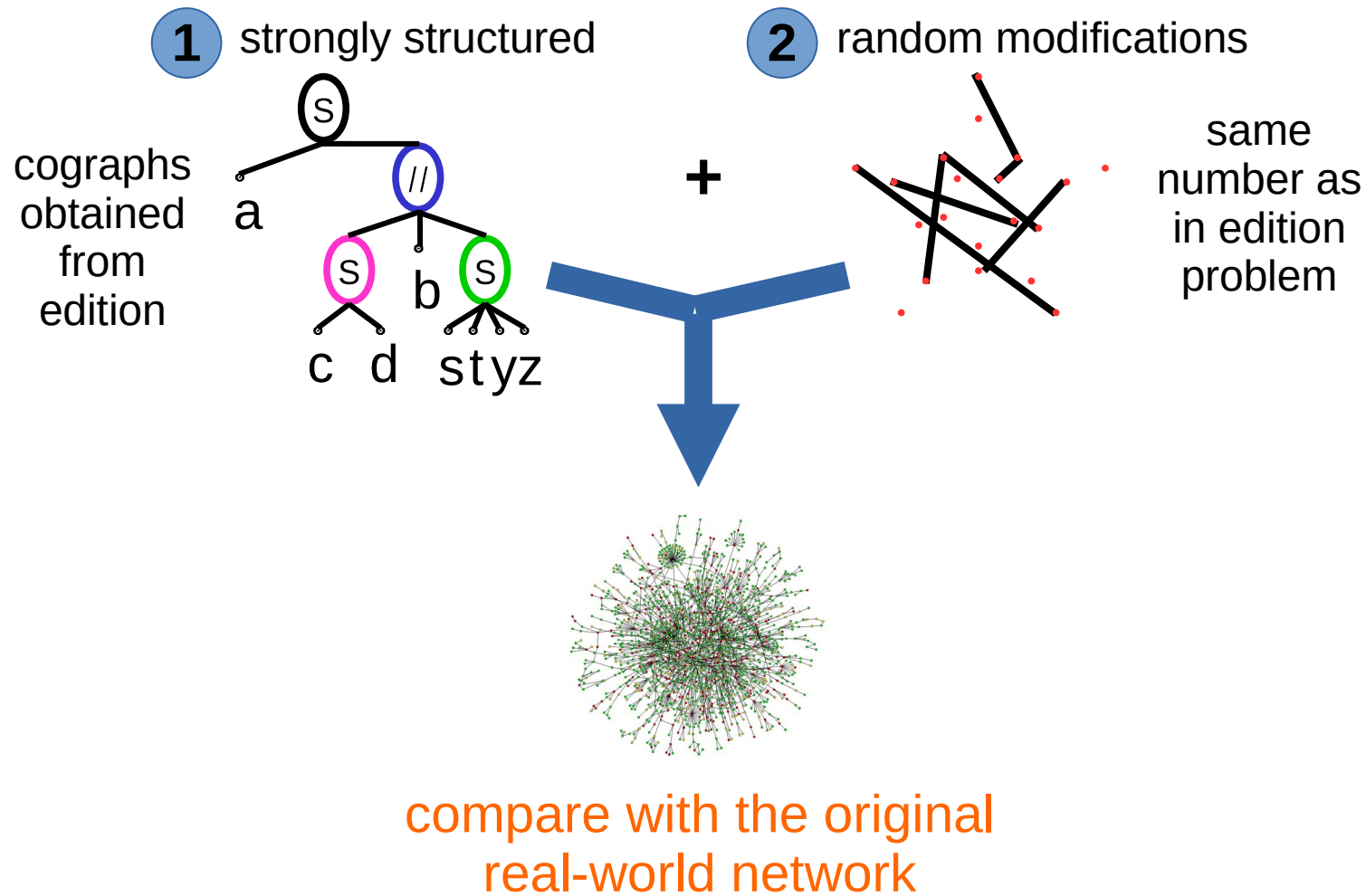
 citation
 social

■ The proximity with cographs highly depends on the real-world context

Testing the modelling approach

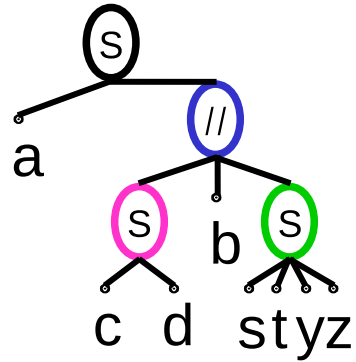


Testing the modelling approach

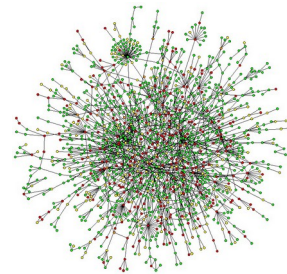
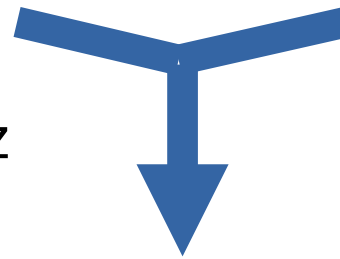
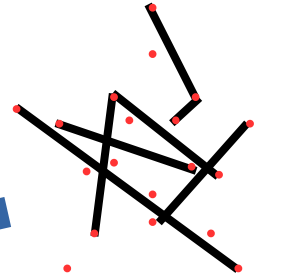


Conclusion

1 strongly structured



2 random modifications

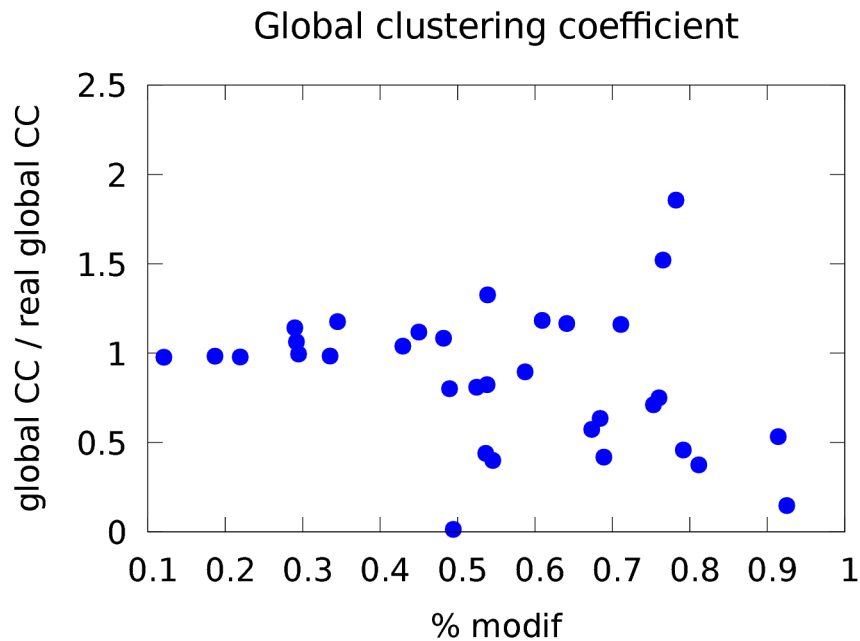


- ✓ global density
- ✓ distances
- ? degree distribution
- ? local density

Results of generation

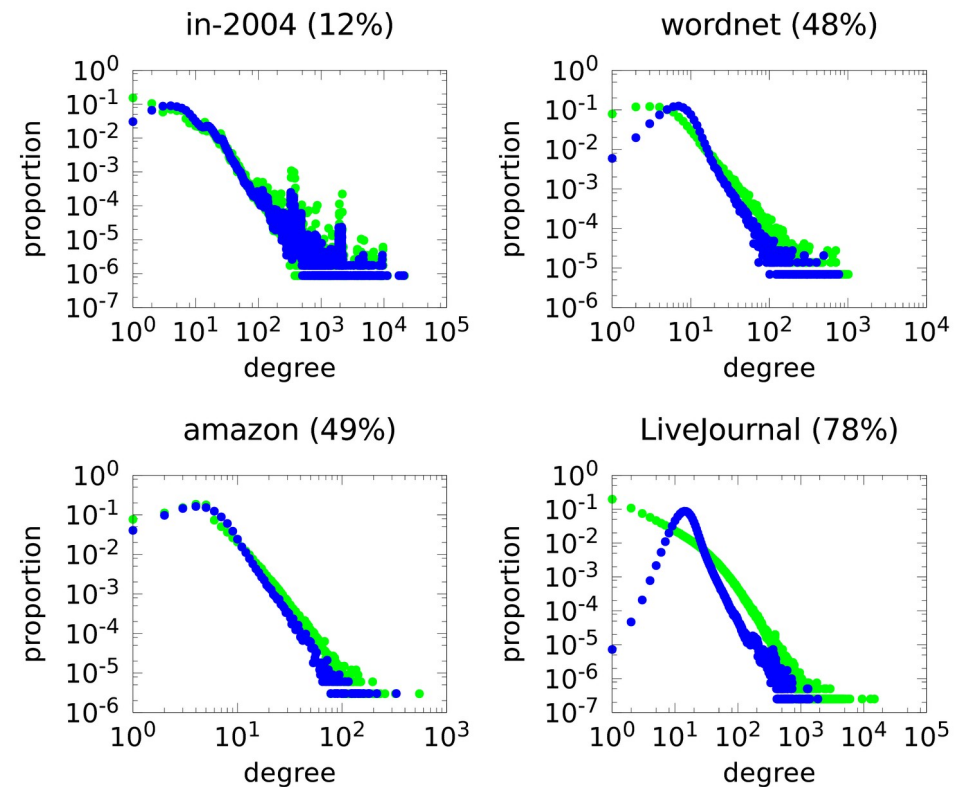
Local density

$$\text{global CC} = \frac{\# \begin{array}{c} \text{?} \\ \text{---} \\ \diagup \quad \diagdown \\ \bullet \quad \bullet \end{array}}{\# \begin{array}{c} \diagup \quad \diagdown \\ \bullet \quad \bullet \end{array}}$$

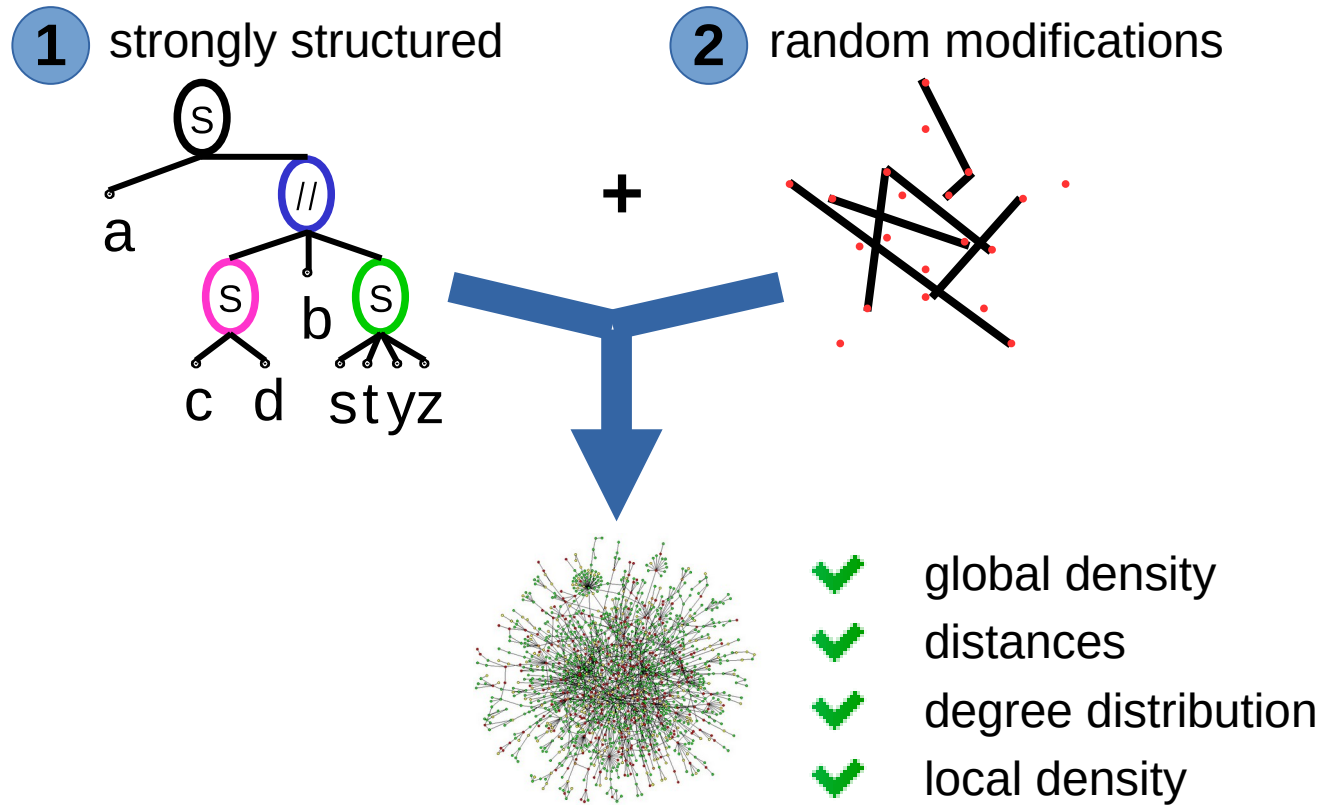


Degree distribution

- Almost cograph model
- Real distribution

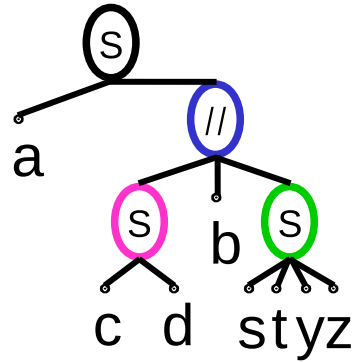


Conclusion

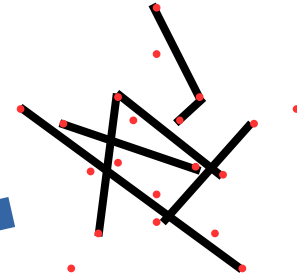


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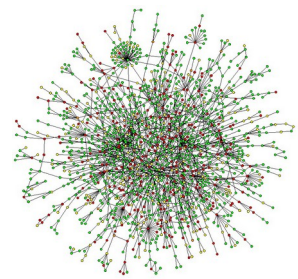
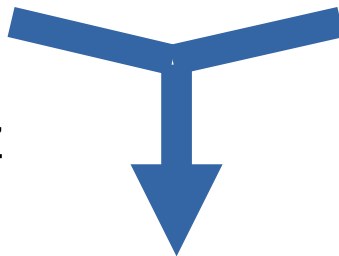
1 strongly structured



2 random modifications



+

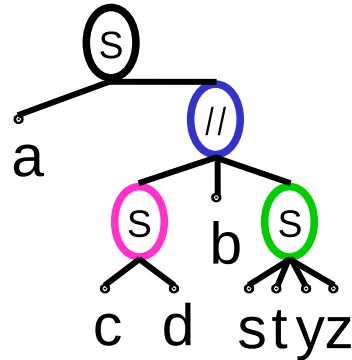


The cograph structure successfully captures these properties

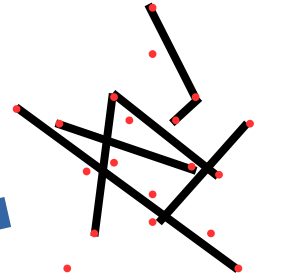
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Conclusion

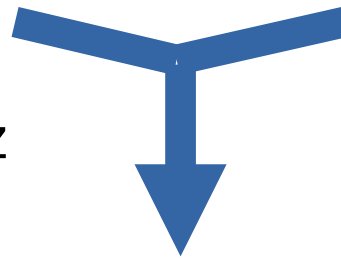
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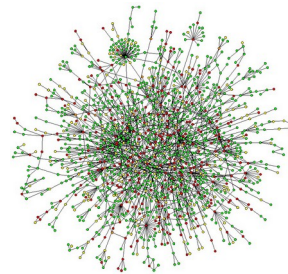
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+



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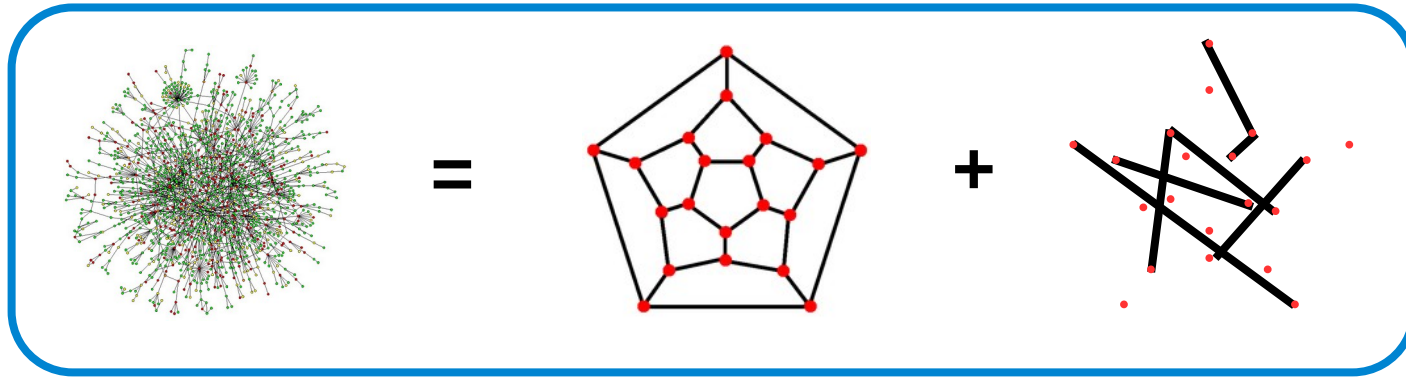
To complete the model

- Edit a real-world graph into a cograph
- ➔ • **Generate a similar cotree**
- Apply random modifications to the cograph

Perspectives

- Complete the modelling approach for cographs
- Consider other graph classes suitable for other kind of networks
 - Chordal graphs → social networks, citations
 - Related to planar graphs → internet, road networks
- Improve algorithms : complexity and quality
 - edition instead of completion
 - avoid incremental approach

Perspectives



- Modelling
- Efficient encoding : space + query time
- Analysis
 - Global organization
 - Specific roles
- Algorithmic theory of ***almost*** structured graphs
 - ➔ Take advantage of the proximity with a strongly structured graph